

# CSC212

# Data Structure

## - Section FG

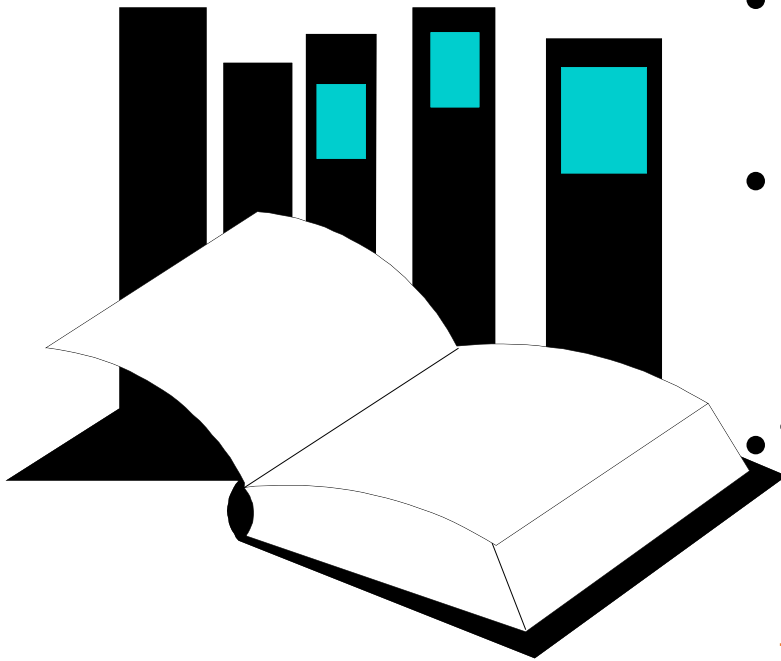
# Lecture 16

# Binary Search Trees

Instructor: Feng HU  
Department of Computer Science  
City College of New York



# Binary Search Trees



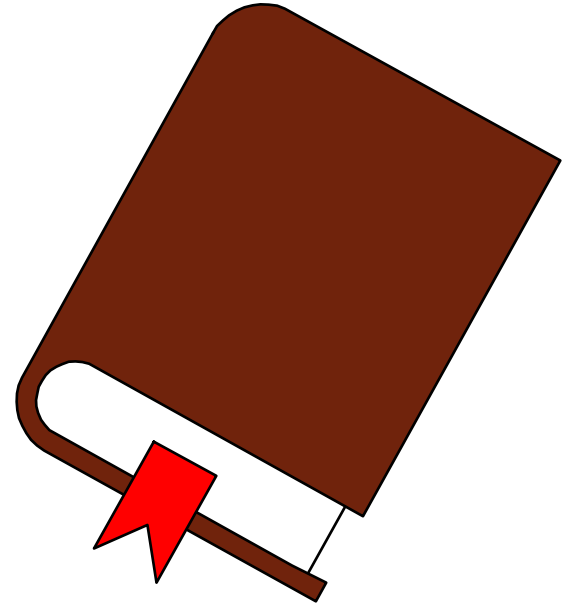
- One of the tree applications in Chapter 10 is binary search trees.
- In Chapter 10, binary search trees are used to implement bags and sets.
- This presentation illustrates how another data type called a dictionary is implemented with binary search trees.

# Binary Search Tree Definition

- In a binary search tree, the entries of the nodes can be compared with a strict weak ordering. Two rules are followed for every node  $n$ :
  - The entry in node  $n$  is NEVER less than an entry in its left subtree
  - The entry in the node  $n$  is less than every entry in its right subtree.

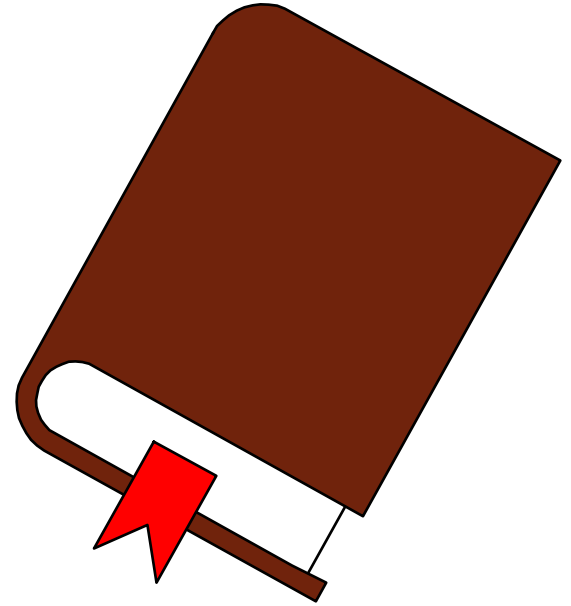
# The Dictionary Data Type

- A dictionary is a collection of items, similar to a bag.
- But unlike a bag, each item has a string attached to it, called the item's key.



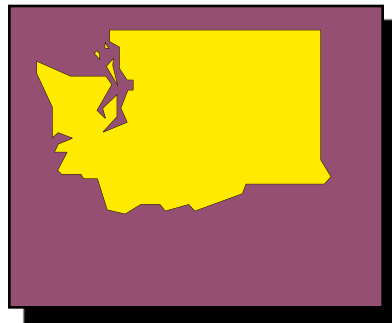
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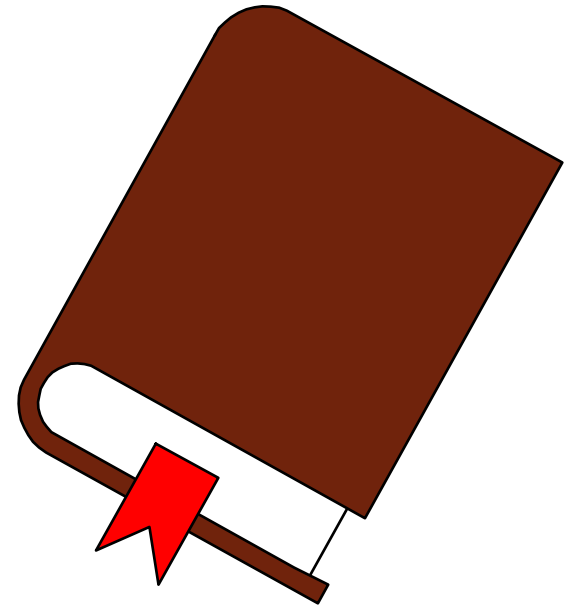
Example:

The items I am storing are records containing data about a state.



# The Dictionary Data Type

- A dictionary is a collection of items, similar to a bag.
- But unlike a bag, each item has a string attached to it, called the item's key.



Example:

The key for each record is the name of the state.



# The Dictionary Data Type

```
void Dictionary::insert(The key for the new item, The new item);
```

- The insertion procedure for a dictionary has two parameters.



# The Dictionary Data Type

- When you want to retrieve an item, you specify the key...

A brown book with a red bookmark is positioned behind a white rectangular box. The book is tilted, and the bookmark is visible at the bottom.

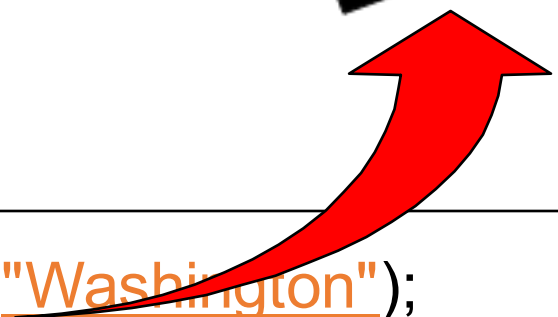
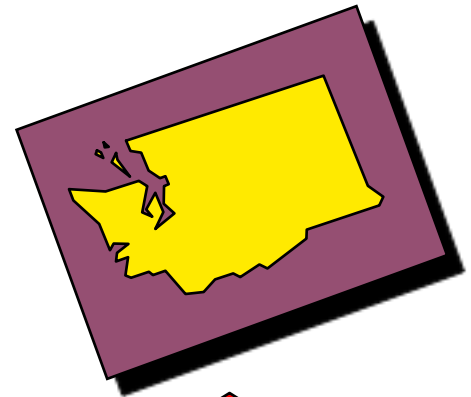
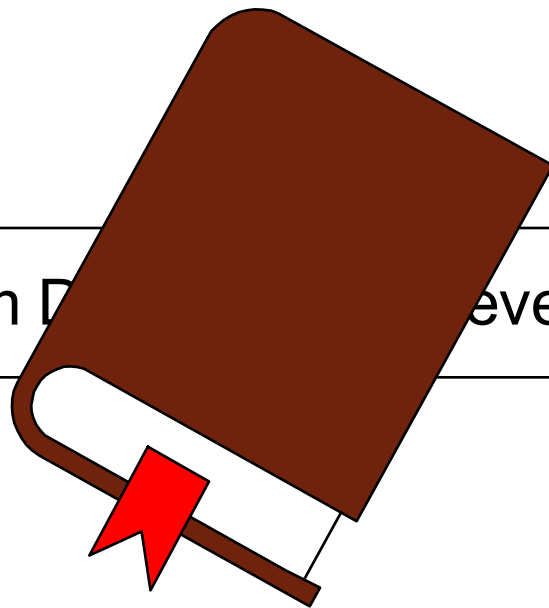
```
Item Dictionary::retrieve("Washington");
```



# The Dictionary Data Type

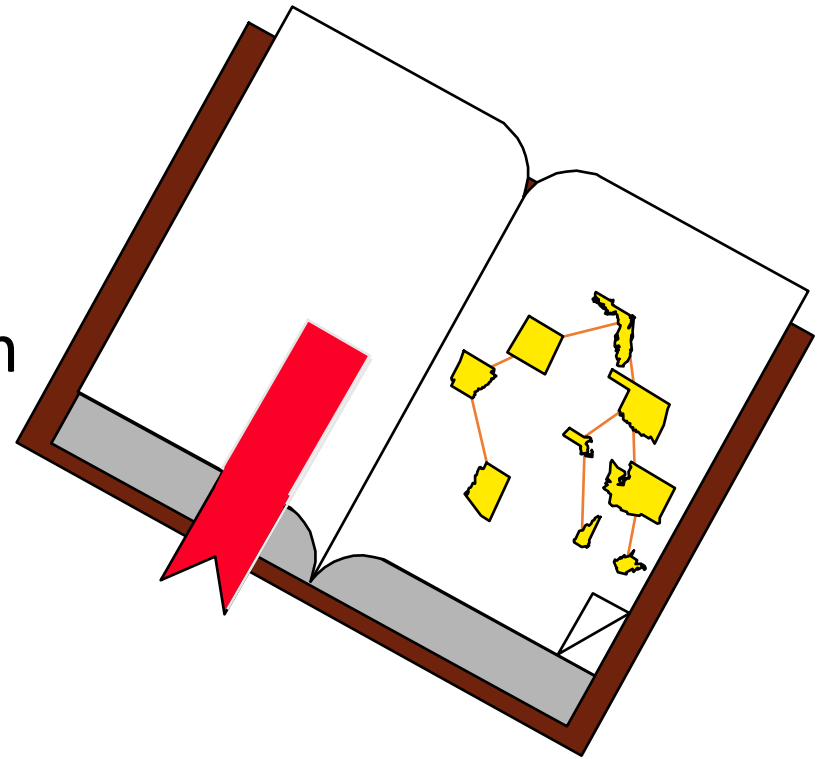
- ❑ When you want to retrieve an item, you specify the key...  
... and the retrieval procedure returns the item.

```
Item D retrieve("Washington");
```



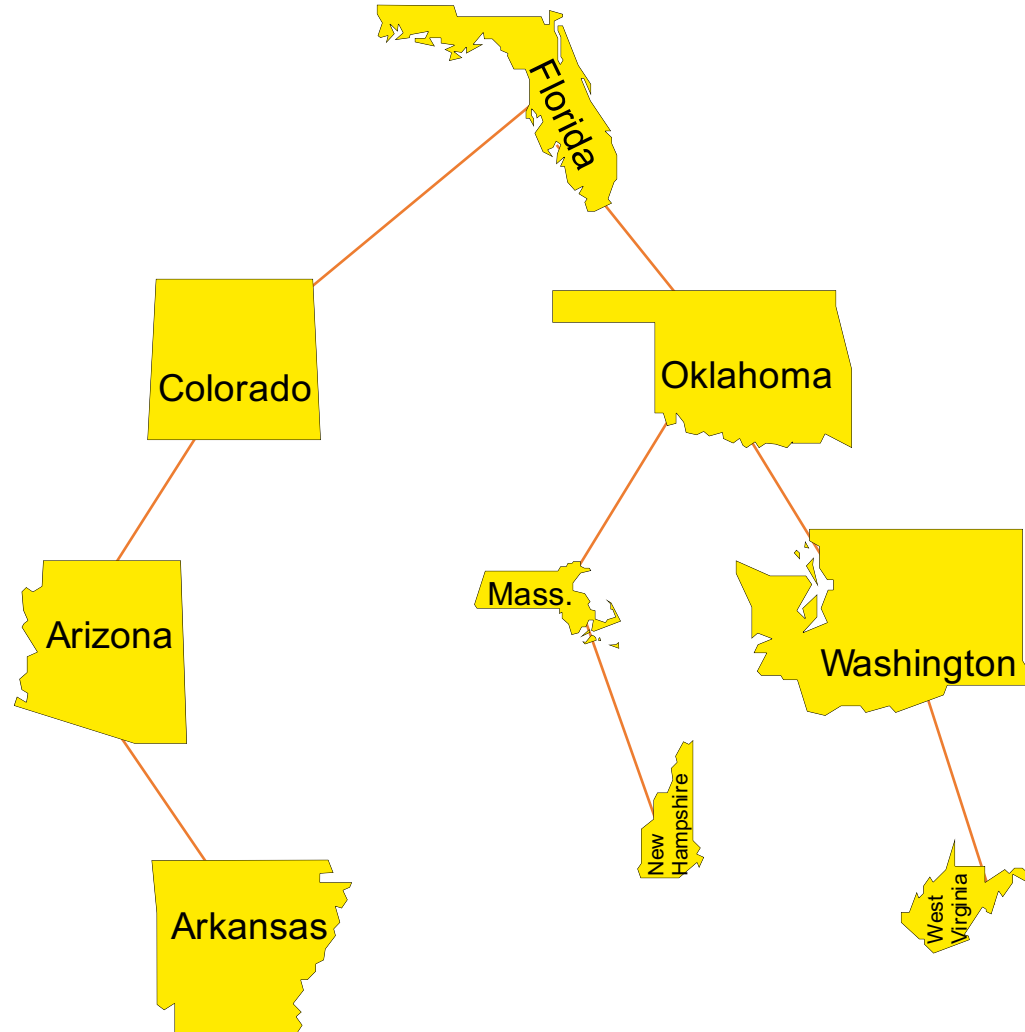
# The Dictionary Data Type

- We'll look at how a binary tree can be used as the internal storage mechanism for the dictionary.



# A Binary Search Tree of States

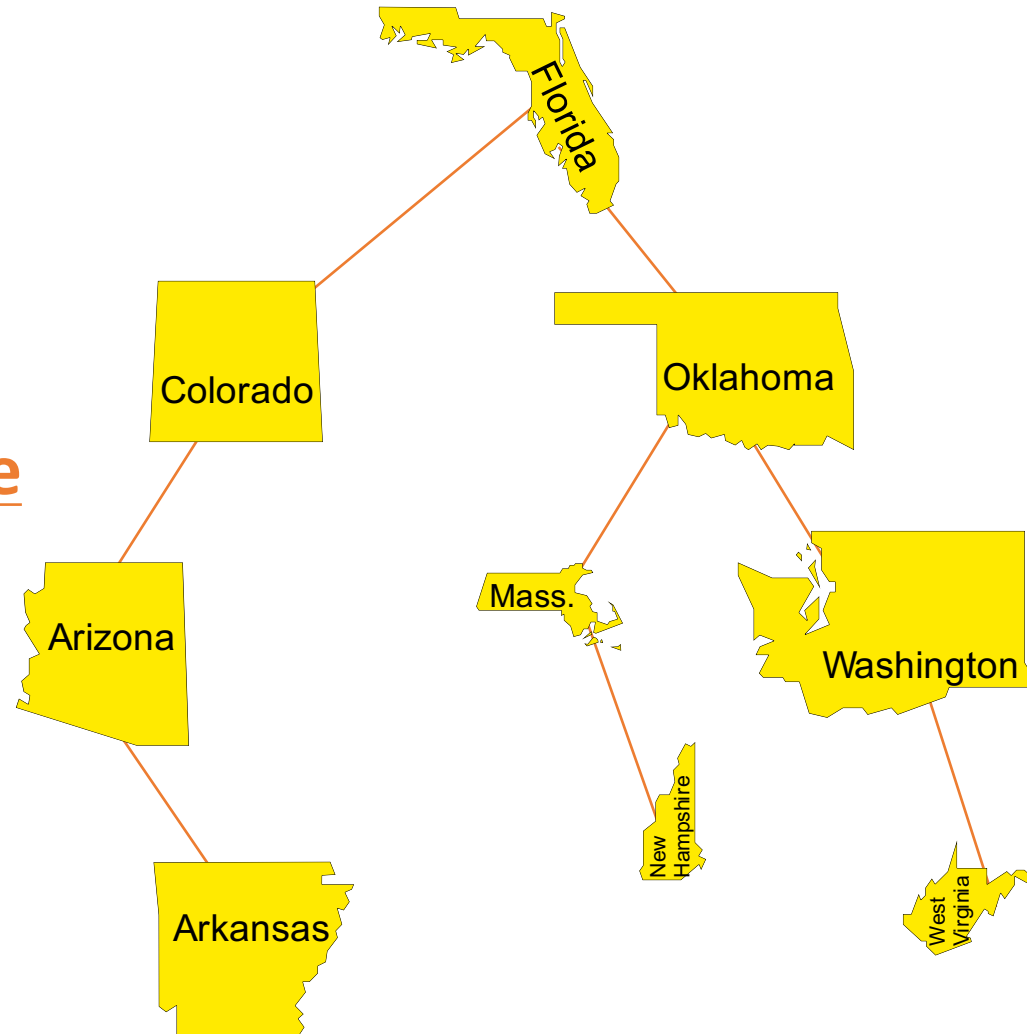
The data in the dictionary will be stored in a binary tree, with each node containing an item and a key.



# A Binary Search Tree of States

Storage rules:

- 1 Every key to the left of a node is alphabetically before the key of the node.



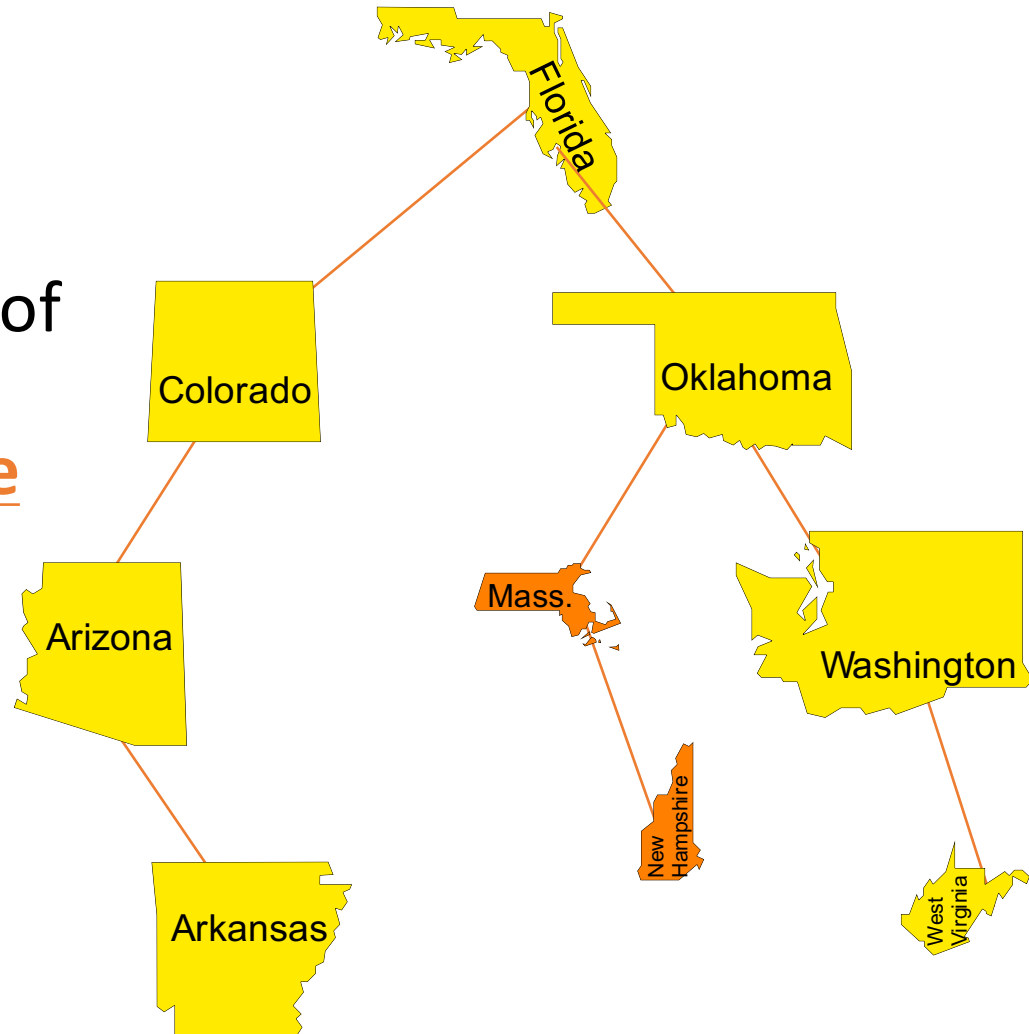
# A Binary Search Tree of States

Storage rules:

- 1 Every key to the left of a node is alphabetically before the key of the node.

Example:

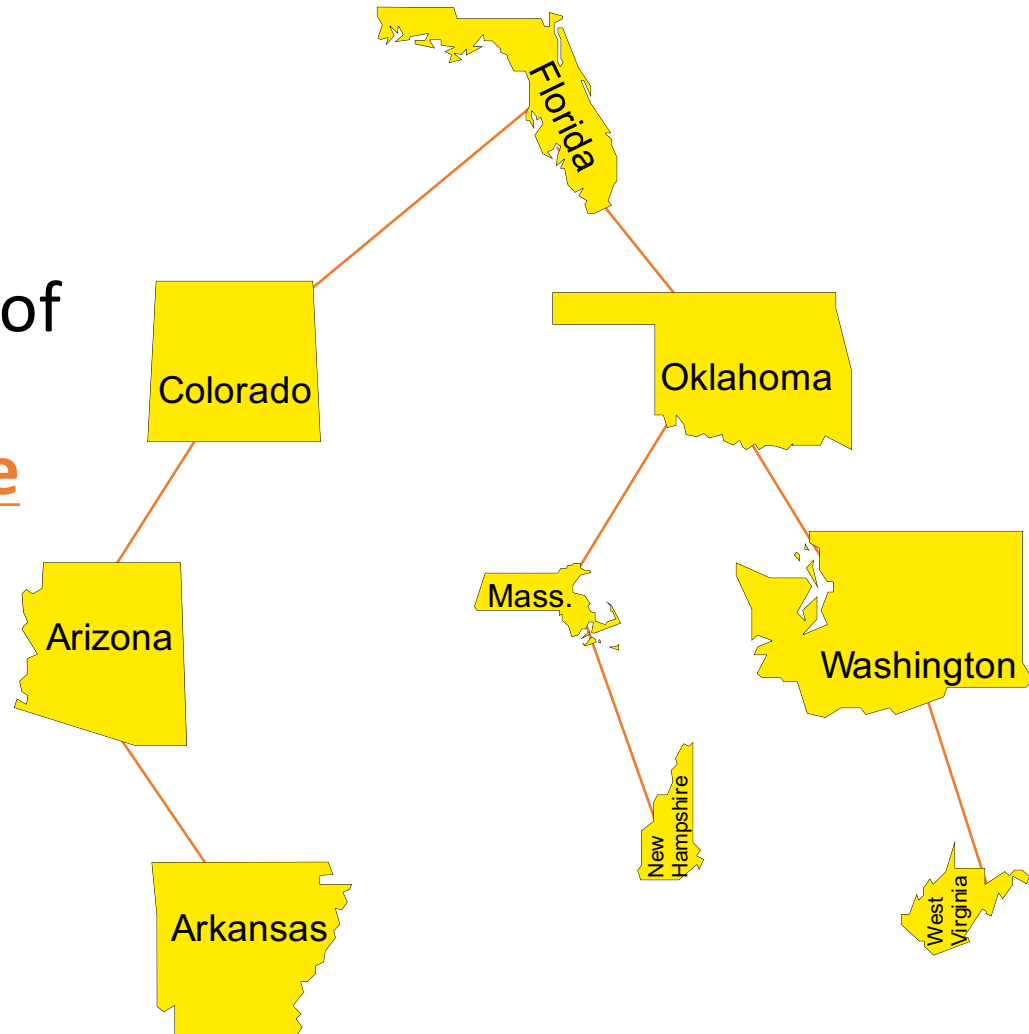
' Massachusetts' and  
' New Hampshire'  
are alphabetically  
before 'Oklahoma'



# A Binary Search Tree of States

Storage rules:

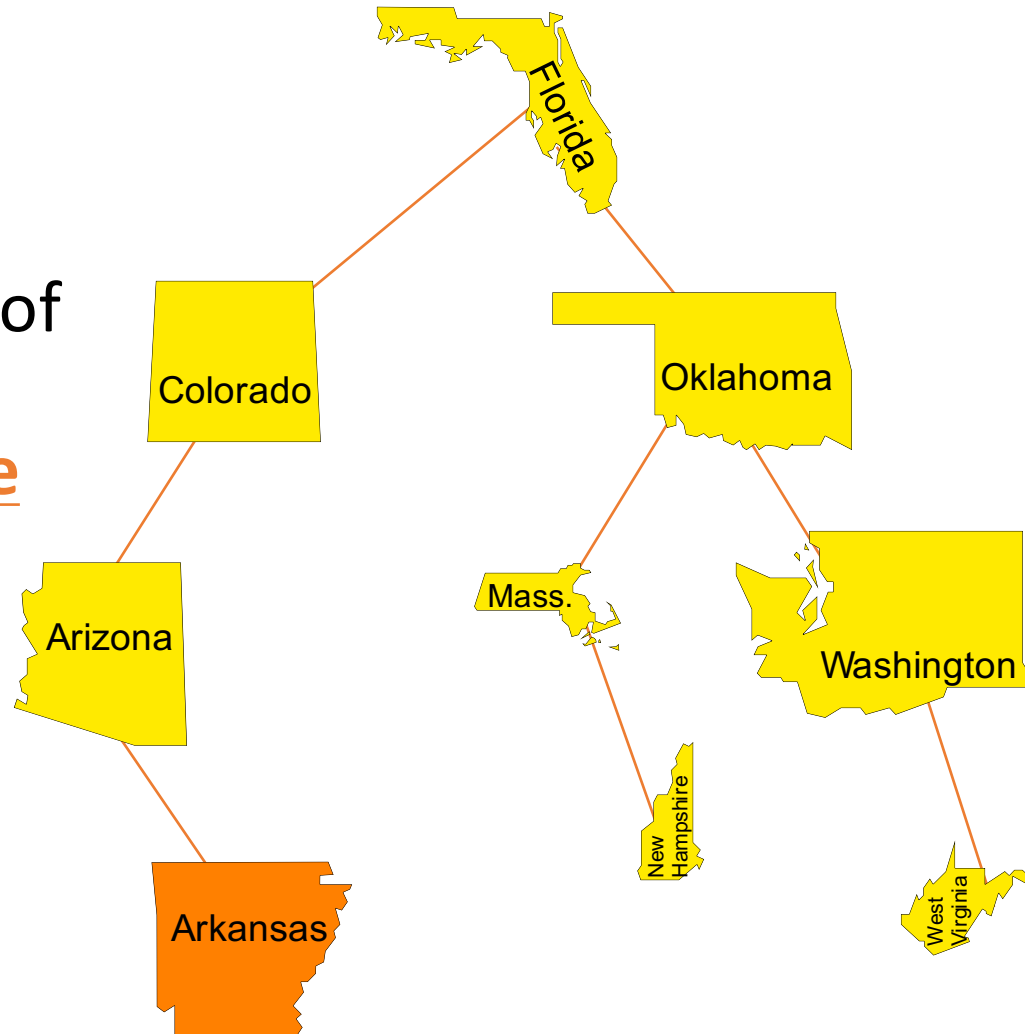
- ① Every key to the left of a node is alphabetically before the key of the node.
- ② Every key to the right of a node is alphabetically after the key of the node.



# A Binary Search Tree of States

Storage rules:

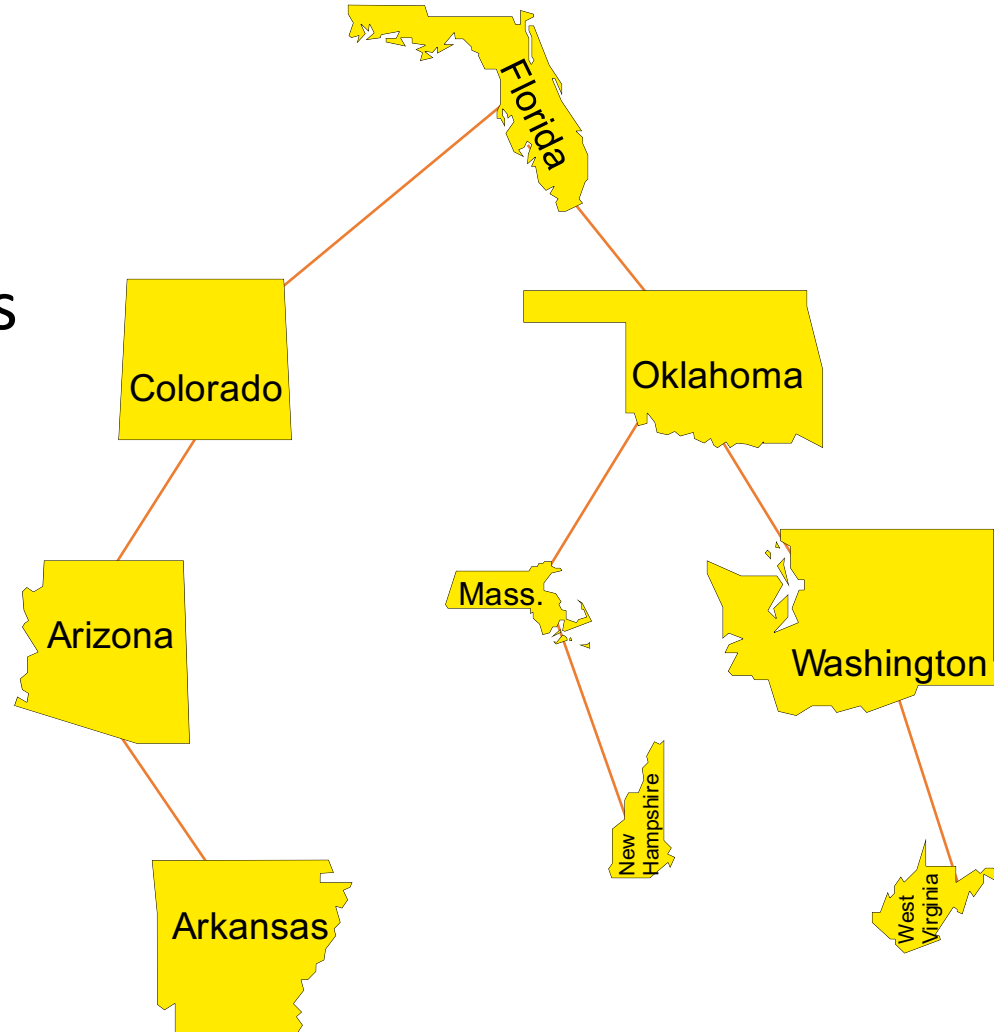
- ① Every key to the left of a node is alphabetically before the key of the node.
- ② Every key to the right of a node is alphabetically after the key of the node.



# Retrieving Data

Start at the root.

- ❶ If the current node has the key, then stop and retrieve the data.
- ❷ If the current node's key is too **large**, move **left** and repeat 1-3.
- ❸ If the current node's key is too **small**, move **right** and repeat 1-3.

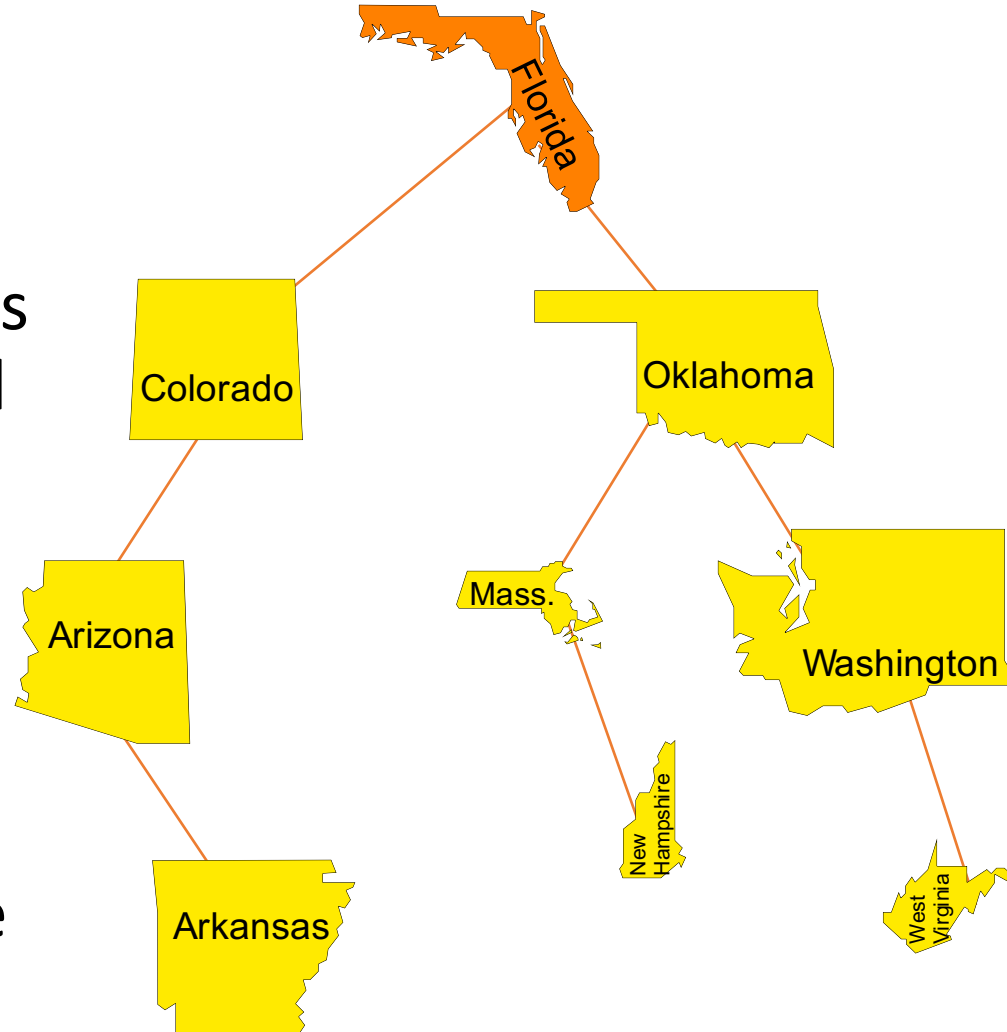




# Retrieve 'New Hampshire'

Start at the root.

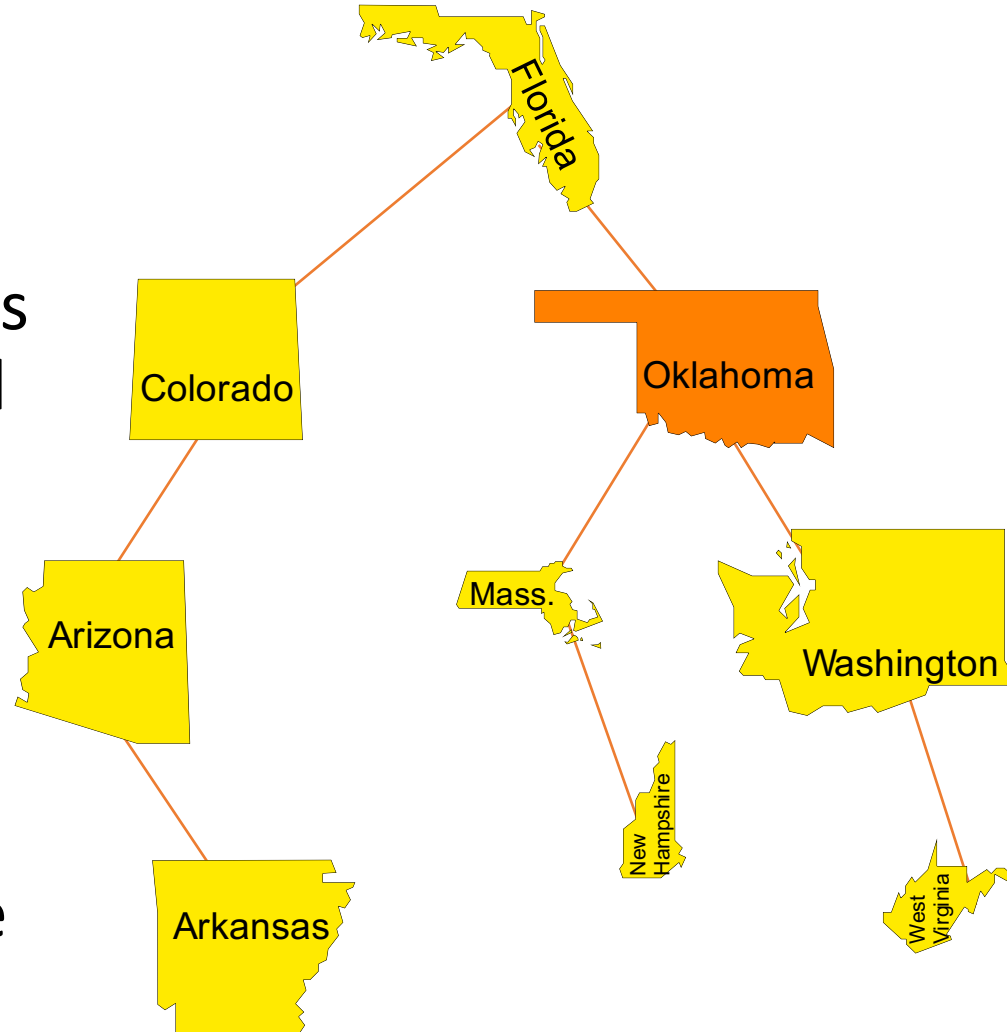
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# Retrieve 'New Hampshire'

Start at the root.

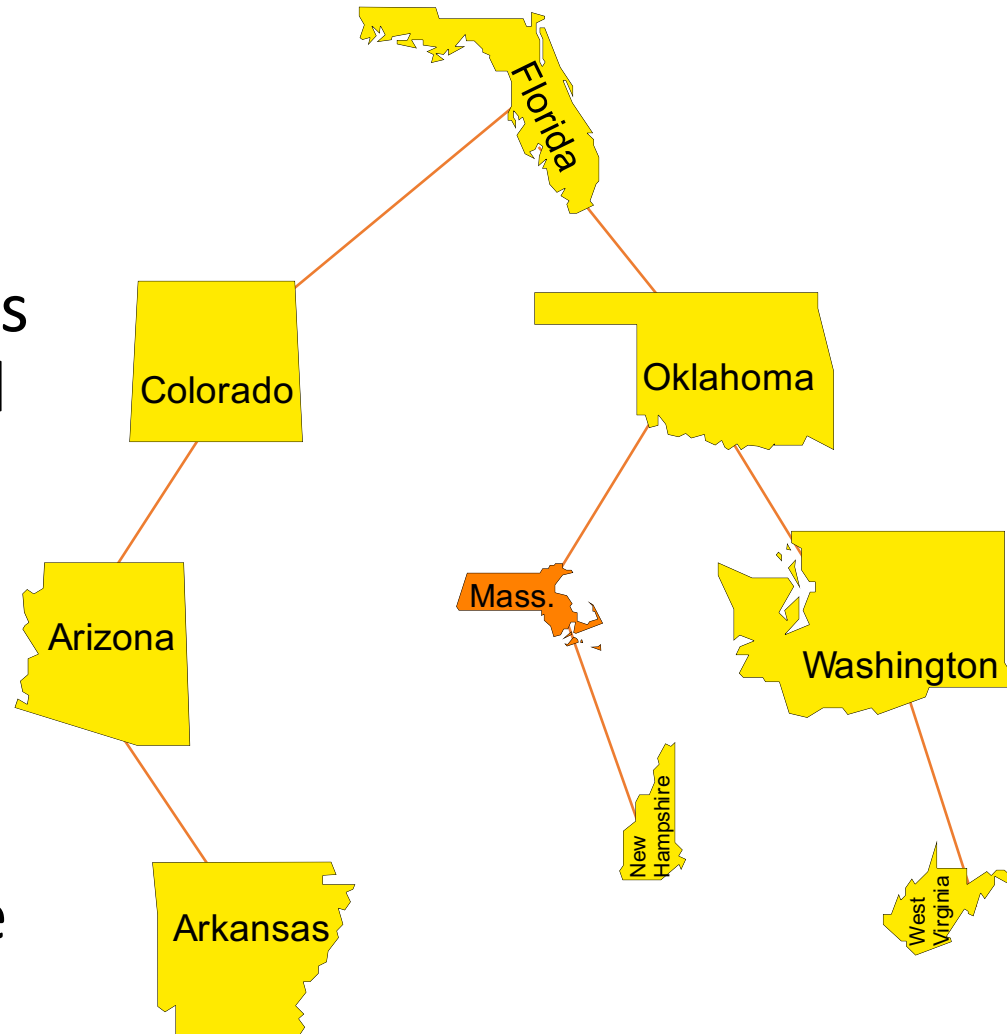
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# Retrieve 'New Hampshire'

Start at the root.

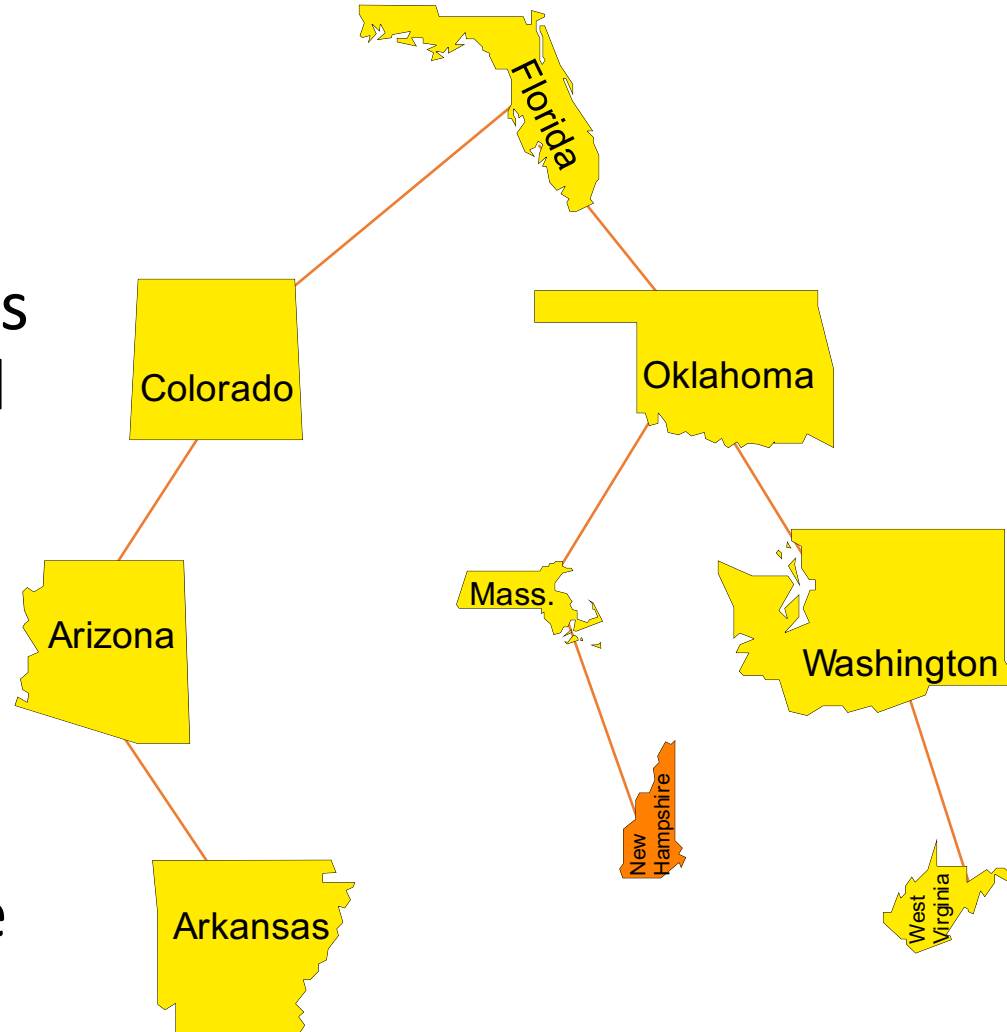
- ❶ If the current node has the key, then stop and retrieve the data.
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# Retrieve 'New Hampshire'

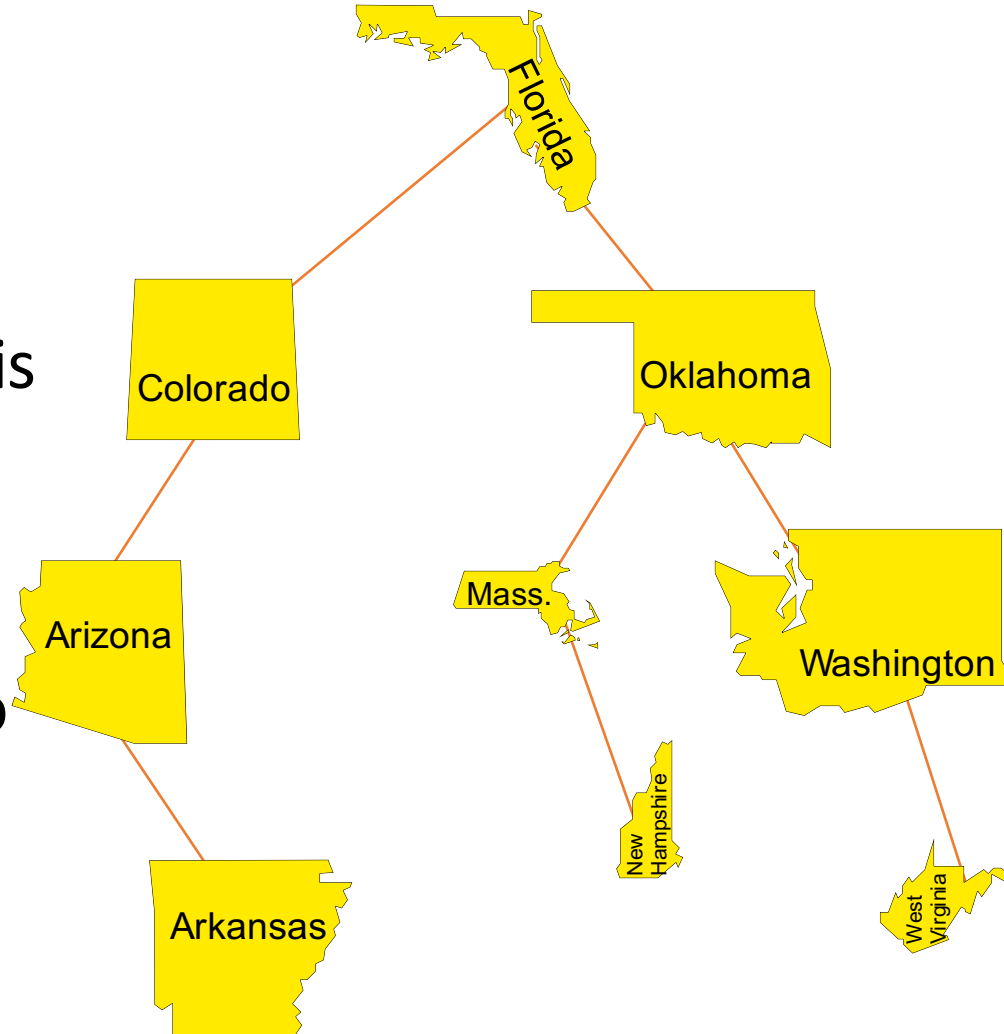
Start at the root.

- ❶ If the current node has the key, then stop and retrieve the data.
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- ❸ If the current node's key is too **small**, move **right** and repeat 1-3.



# Adding a New Item with a Given Key

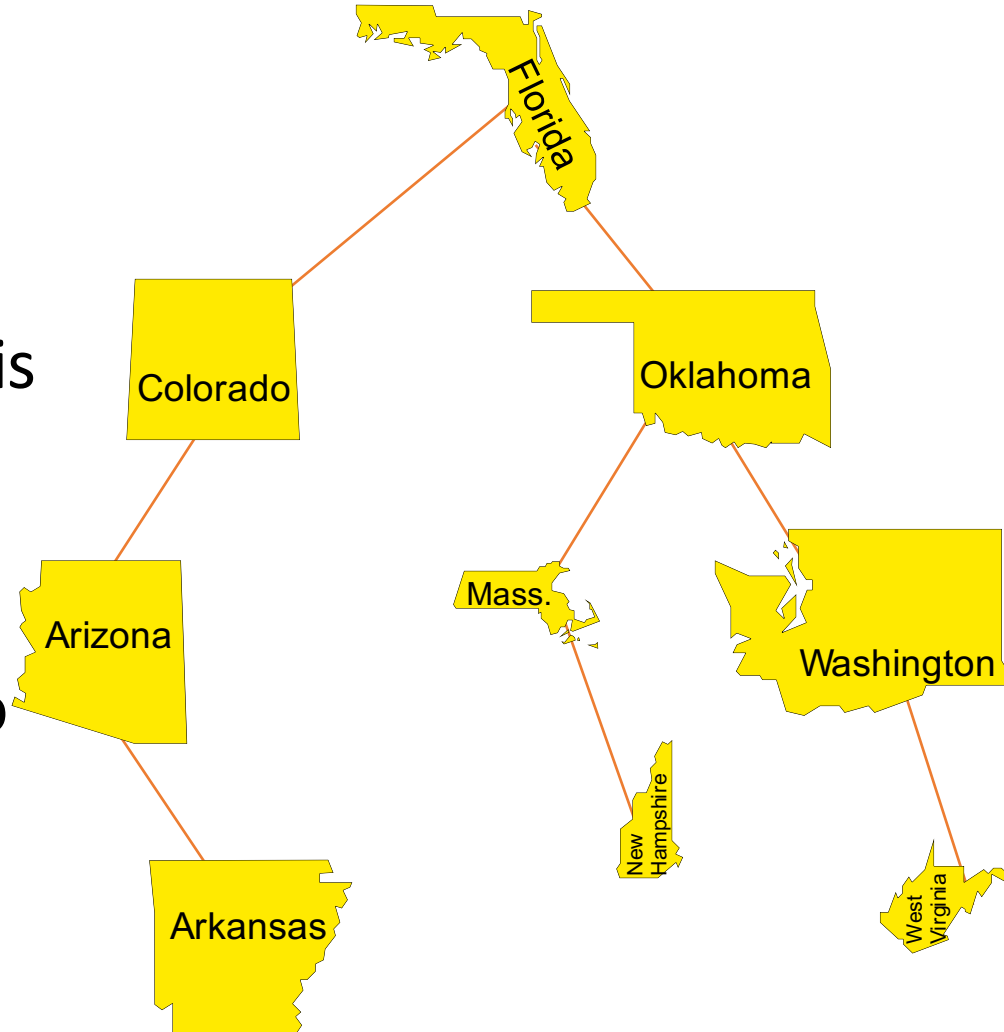
- ❶ Pretend that you are trying to find the key, but stop when there is no node to move to.
- ❷ Add the new node at the spot where you would have moved to if there had been a node.



# Adding



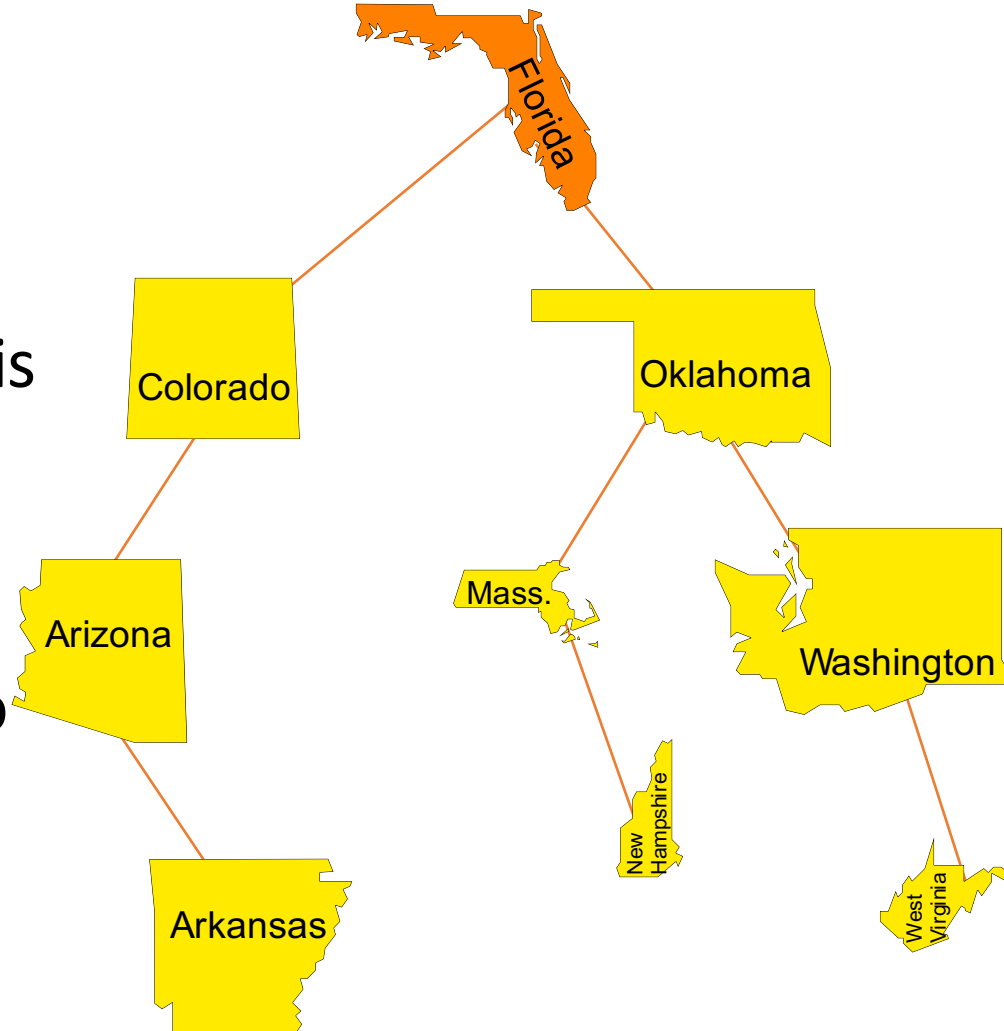
- 1 Pretend that you are trying to find the key, but stop when there is no node to move to.
- 2 Add the new node at the spot where you would have moved to if there had been a node.



# Adding



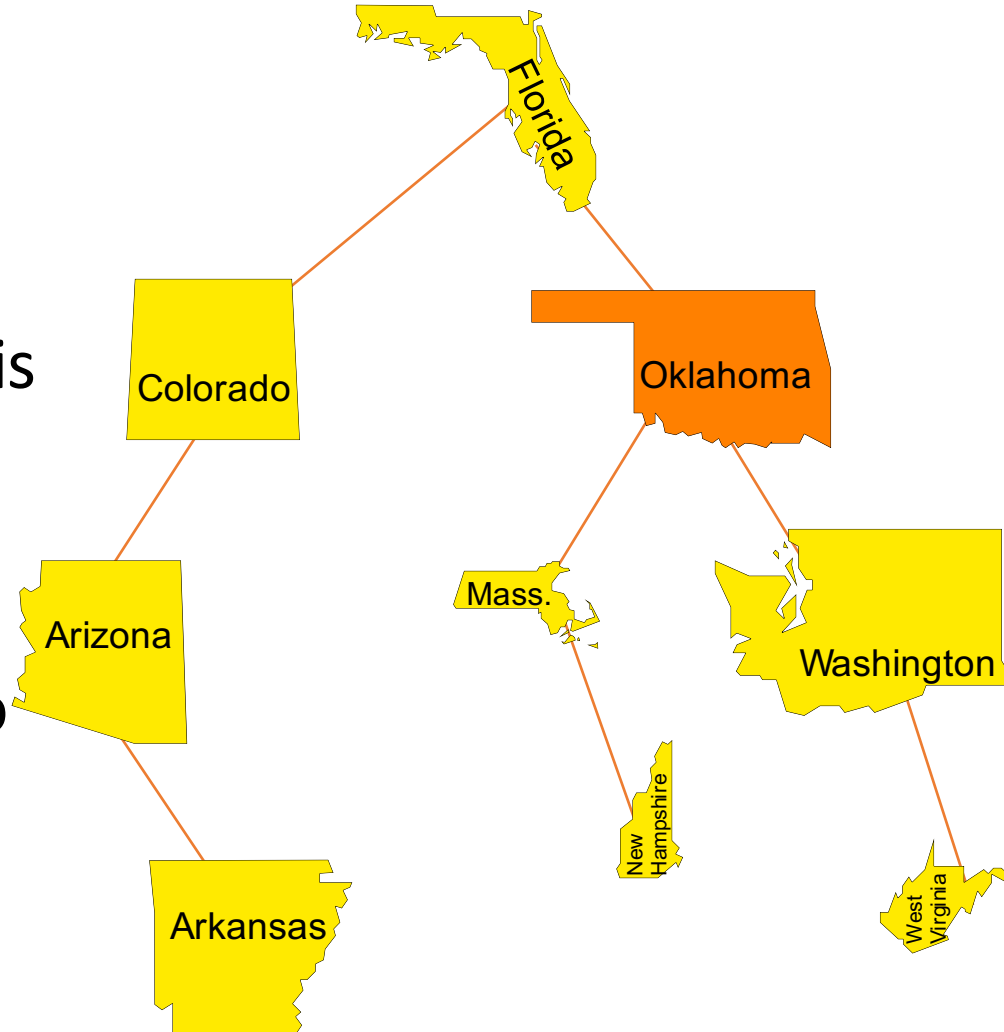
- 1 Pretend that you are trying to find the key, but stop when there is no node to move to.
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# Adding



- 1 Pretend that you are trying to find the key, but stop when there is no node to move to.
- 2 Add the new node at the spot where you would have moved to if there had been a node.

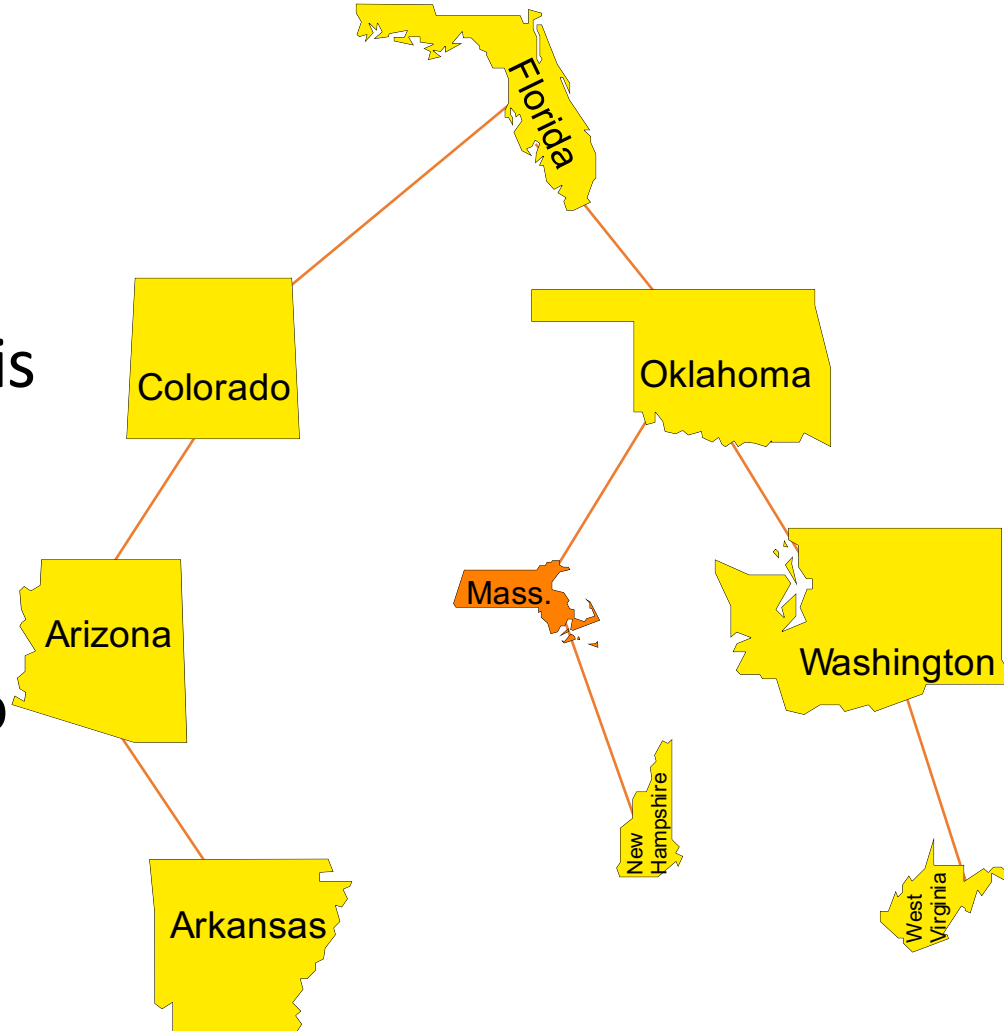




# Adding



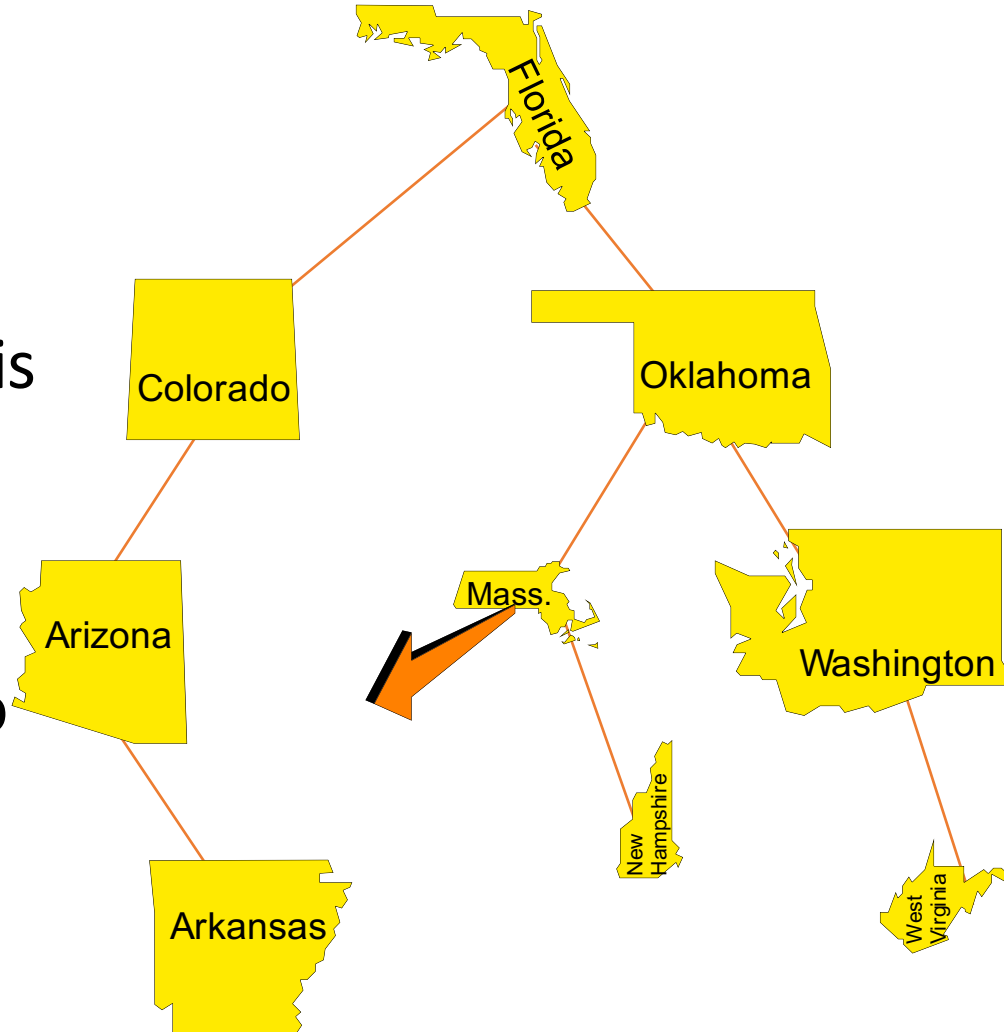
- 1 Pretend that you are trying to find the key, but stop when there is no node to move to.
- 2 Add the new node at the spot where you would have moved to if there had been a node.



# Adding

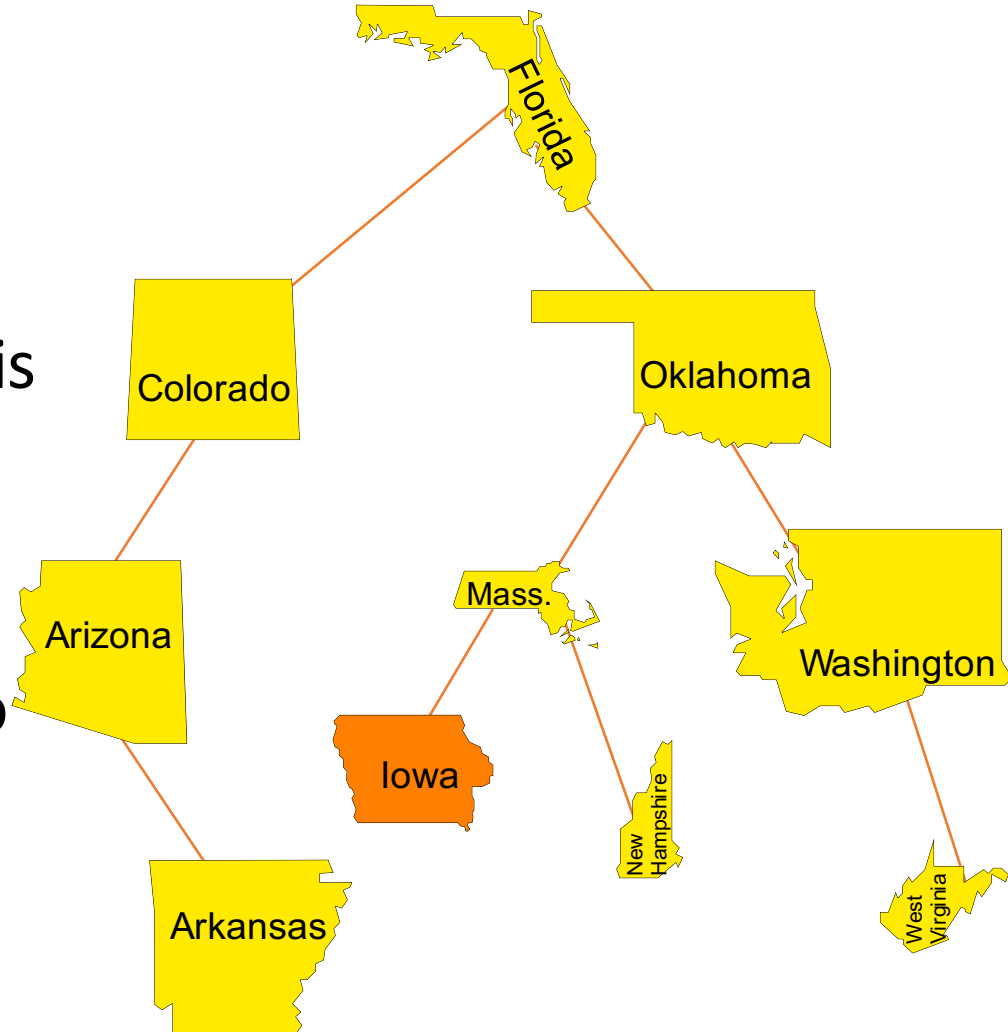


- 1 Pretend that you are trying to find the key, but stop when there is no node to move to.
- 2 Add the new node at the spot where you would have moved to if there had been a node.



# Adding

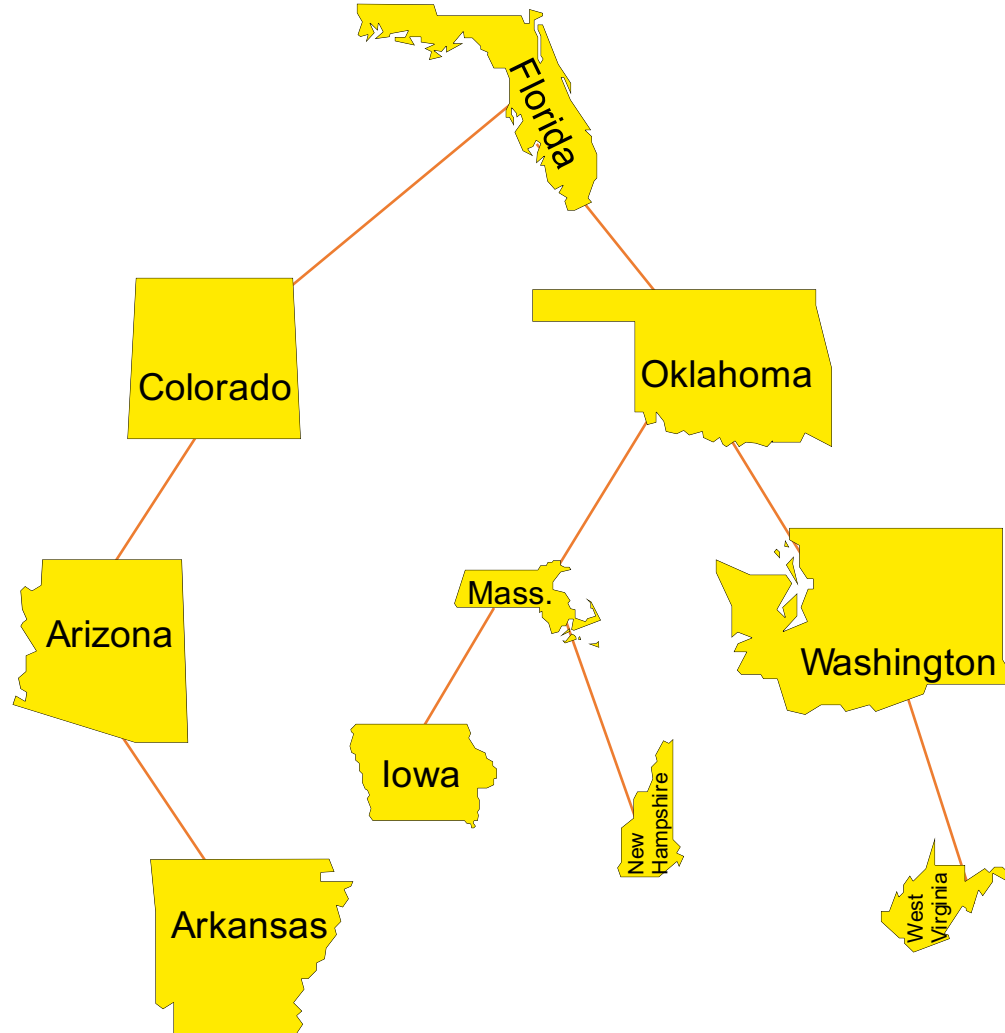
- ❶ Pretend that you are trying to find the key, but stop when there is no node to move to.
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Adding

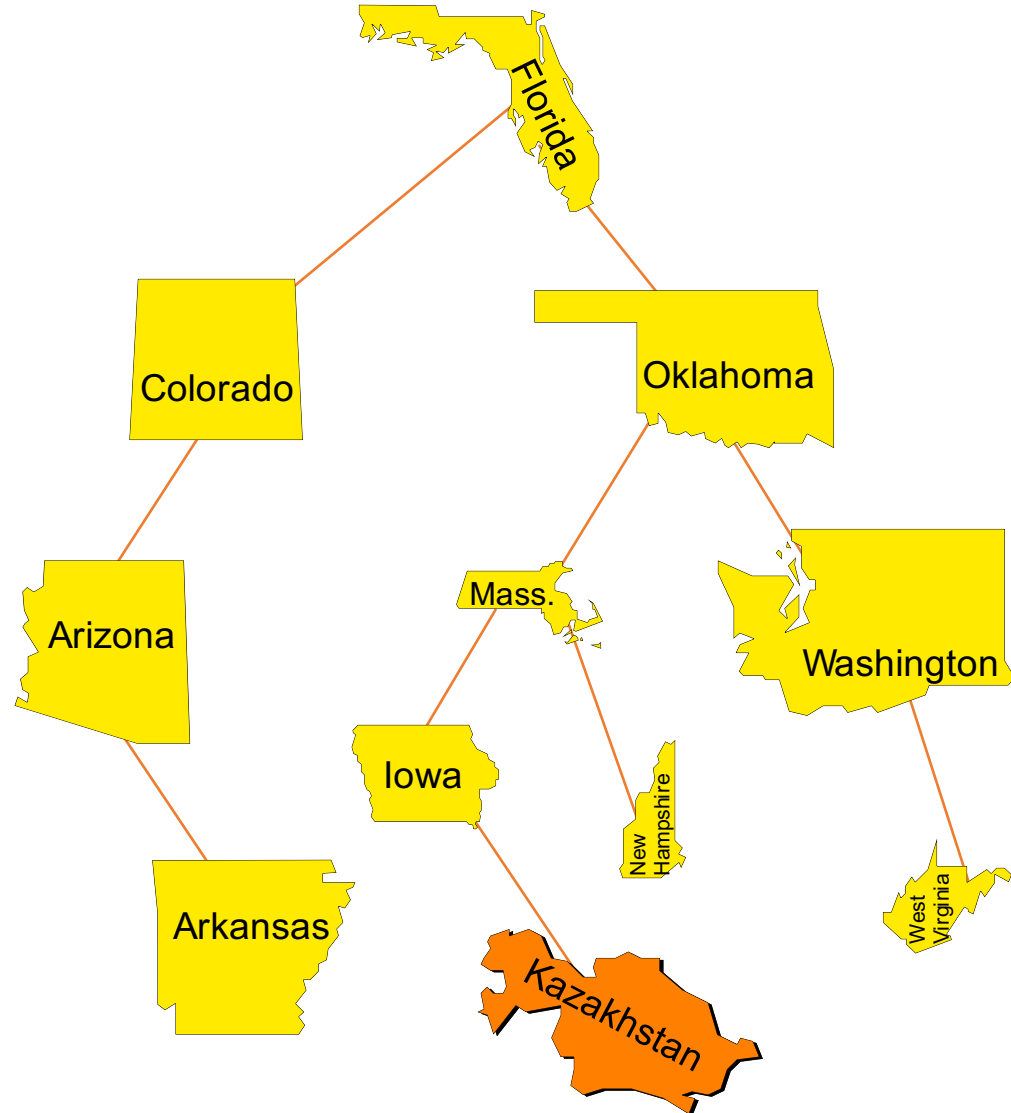


*Where would you add this state?*



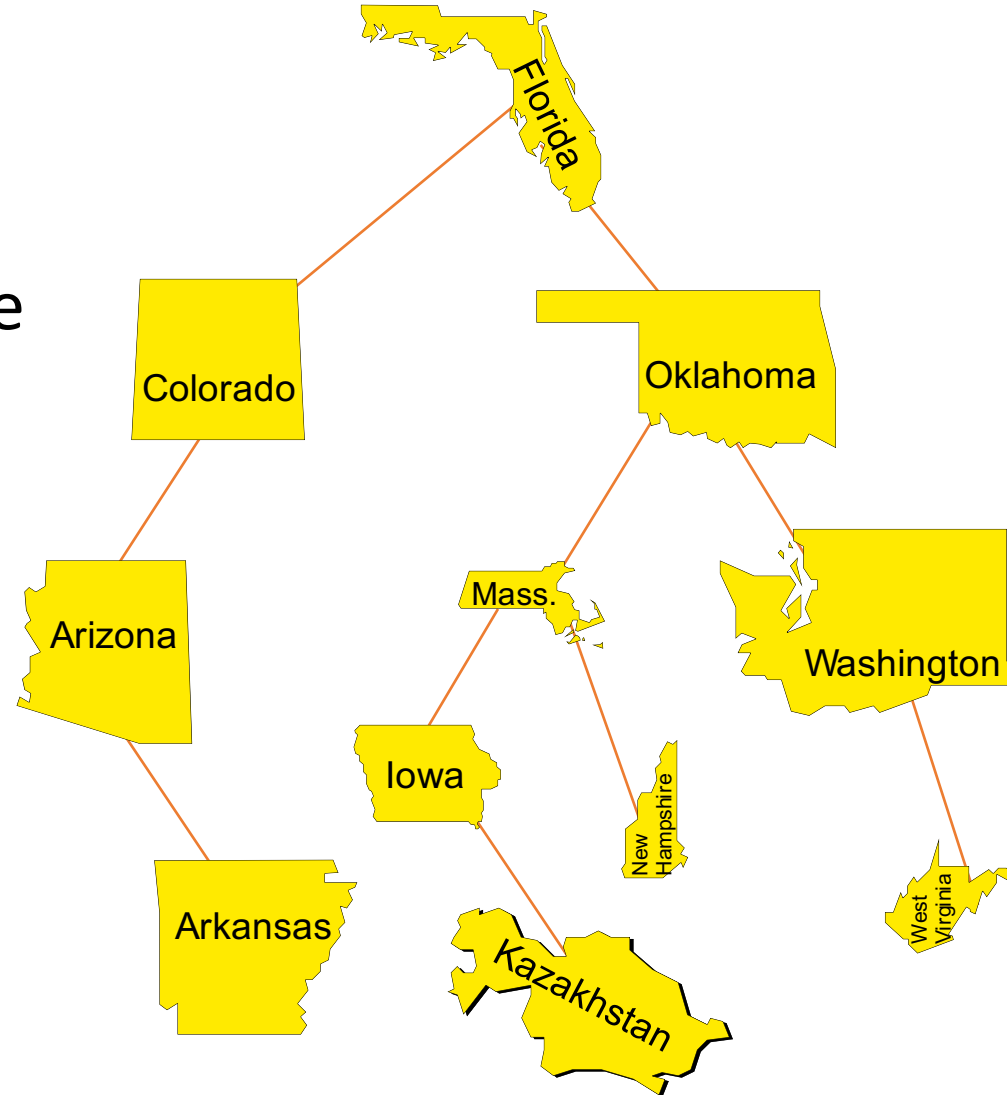
# Adding

Kazakhstan is the  
new right child  
of Iowa?



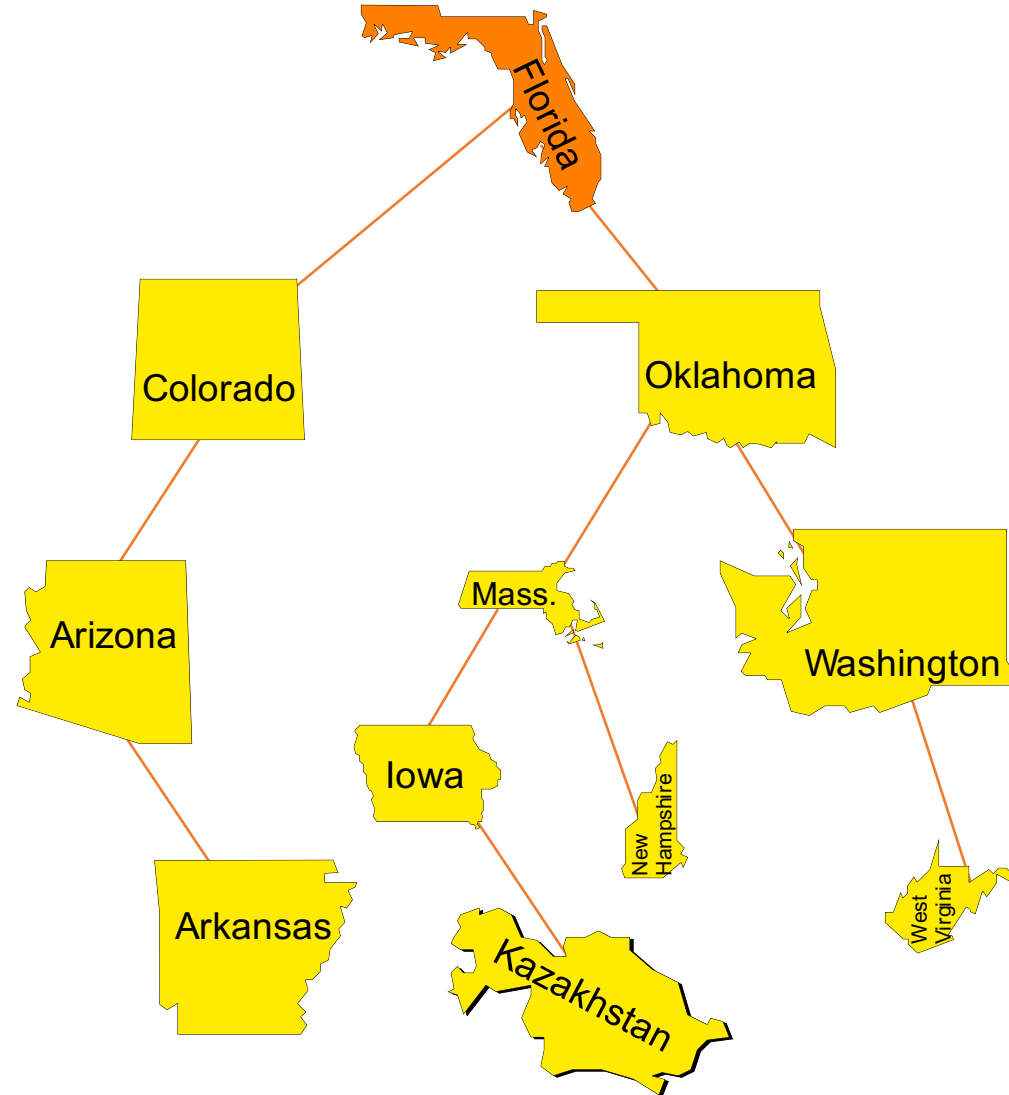
# Removing an Item with a Given Key

- 1 Find the item.
- 2 If necessary, swap the item with one that is easier to remove.
- 3 Remove the item.



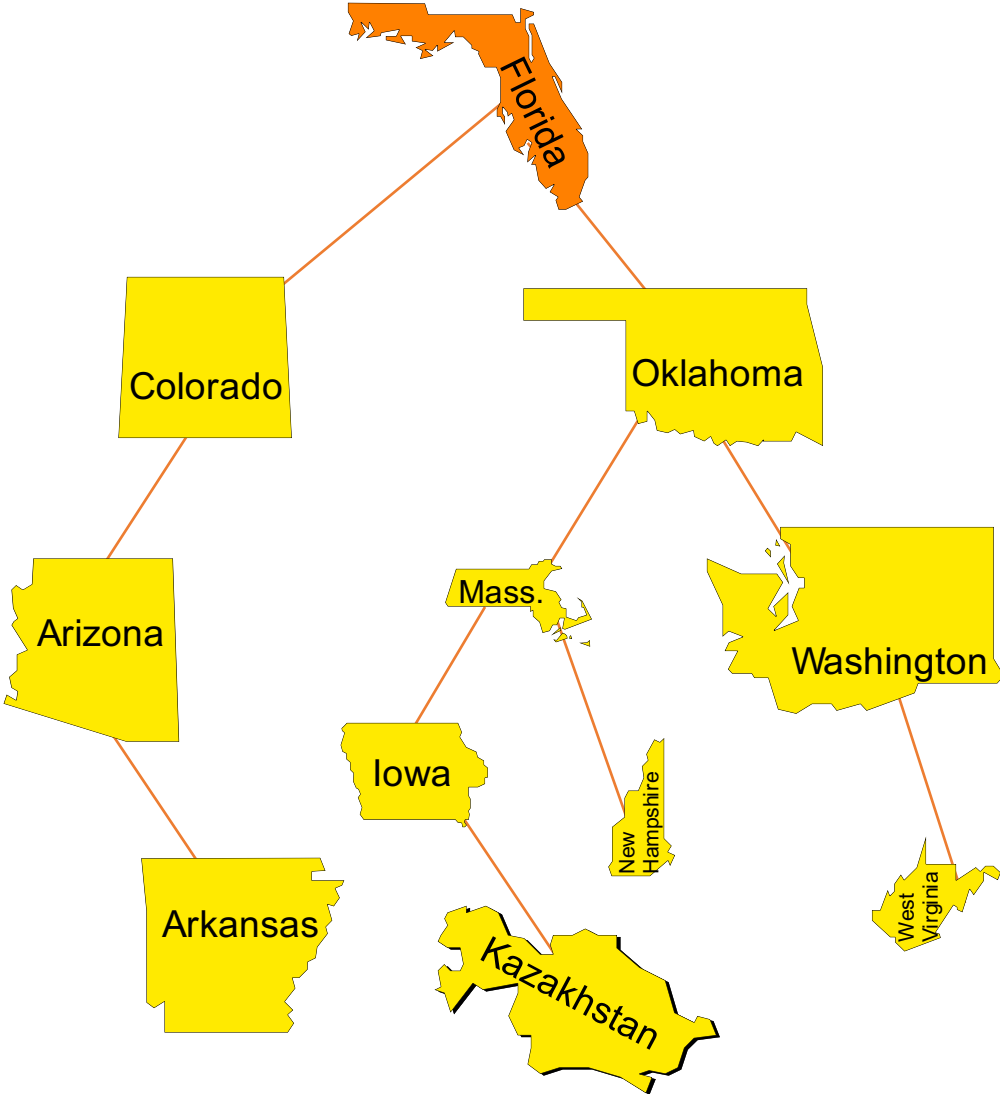
# Removing 'Florida'

① Find the item.



# Removing 'Florida'

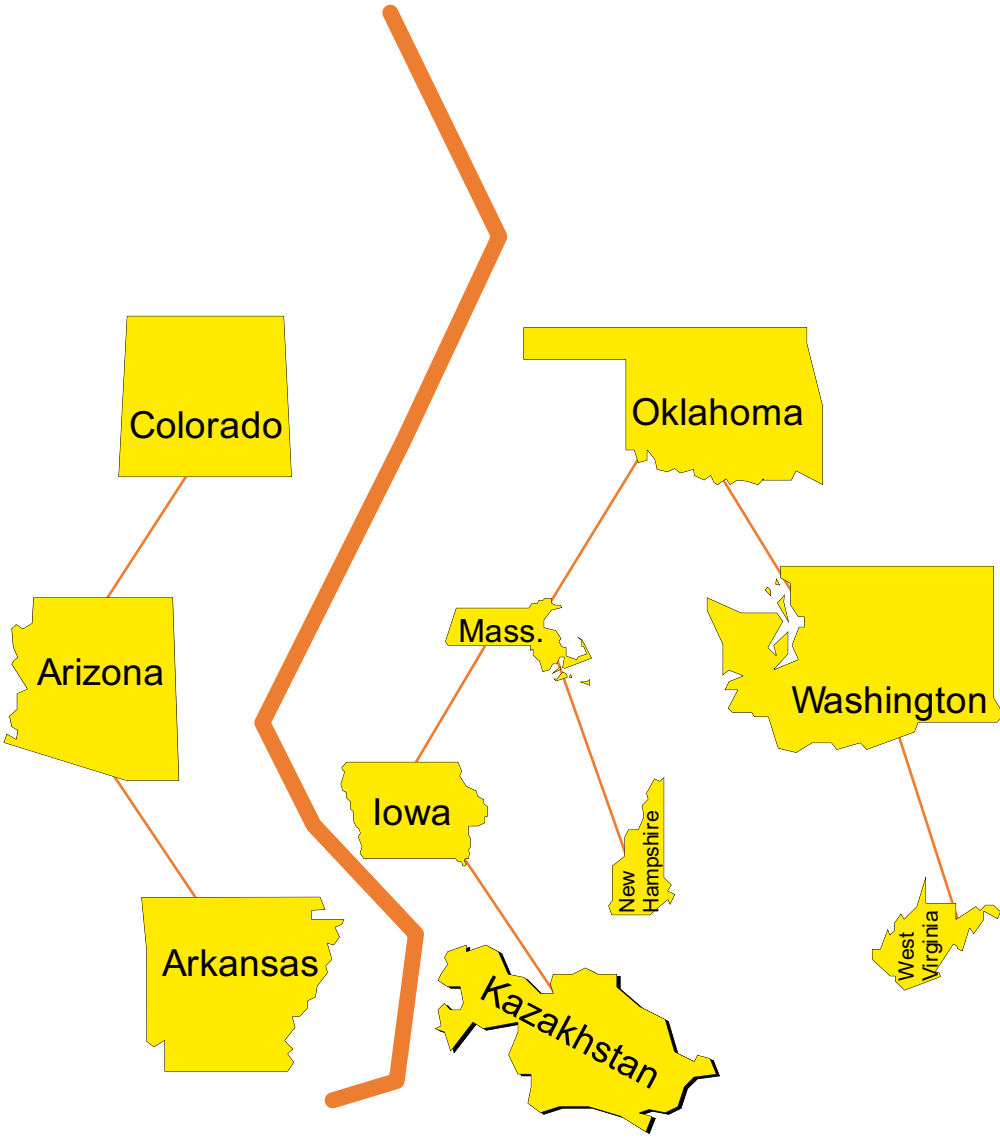
Florida cannot be removed at the moment...





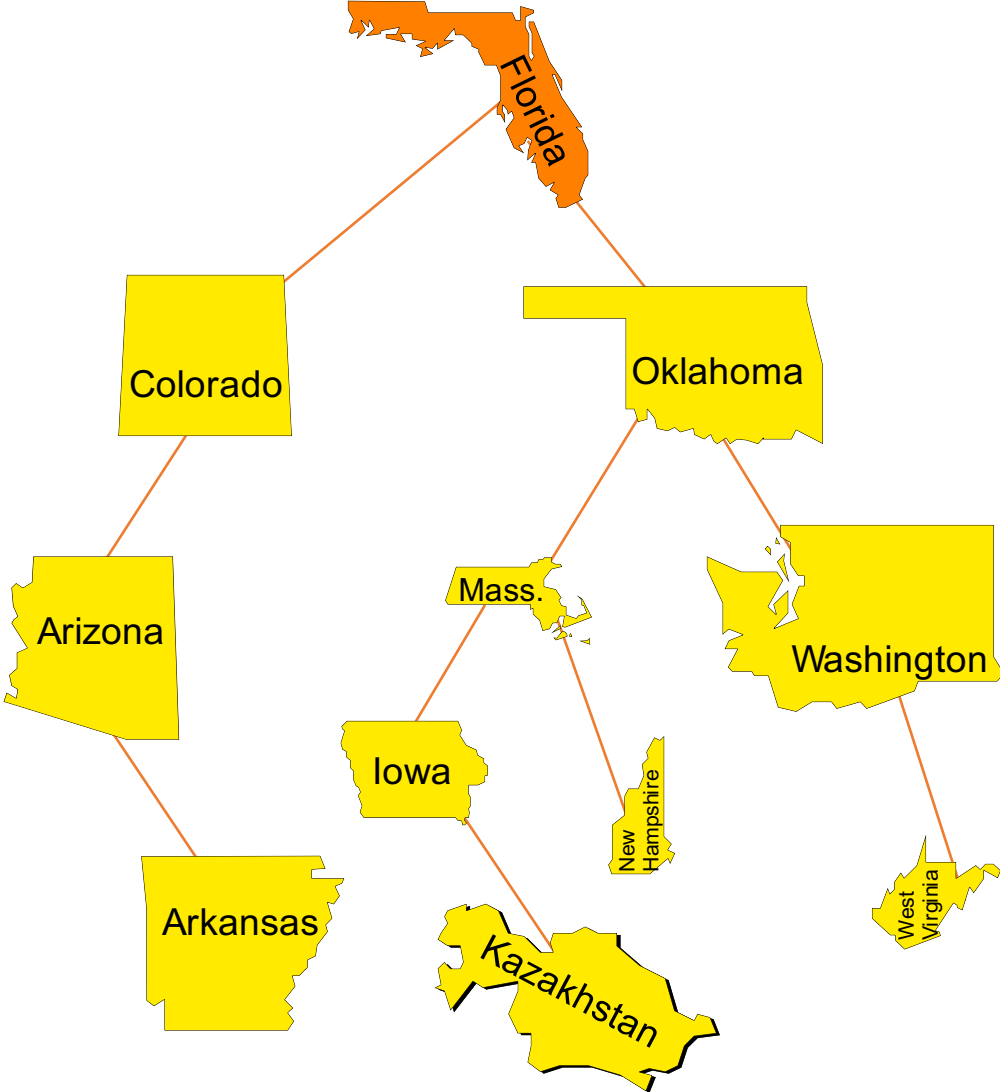
# Removing 'Florida'

... because removing Florida would break the tree into two pieces.



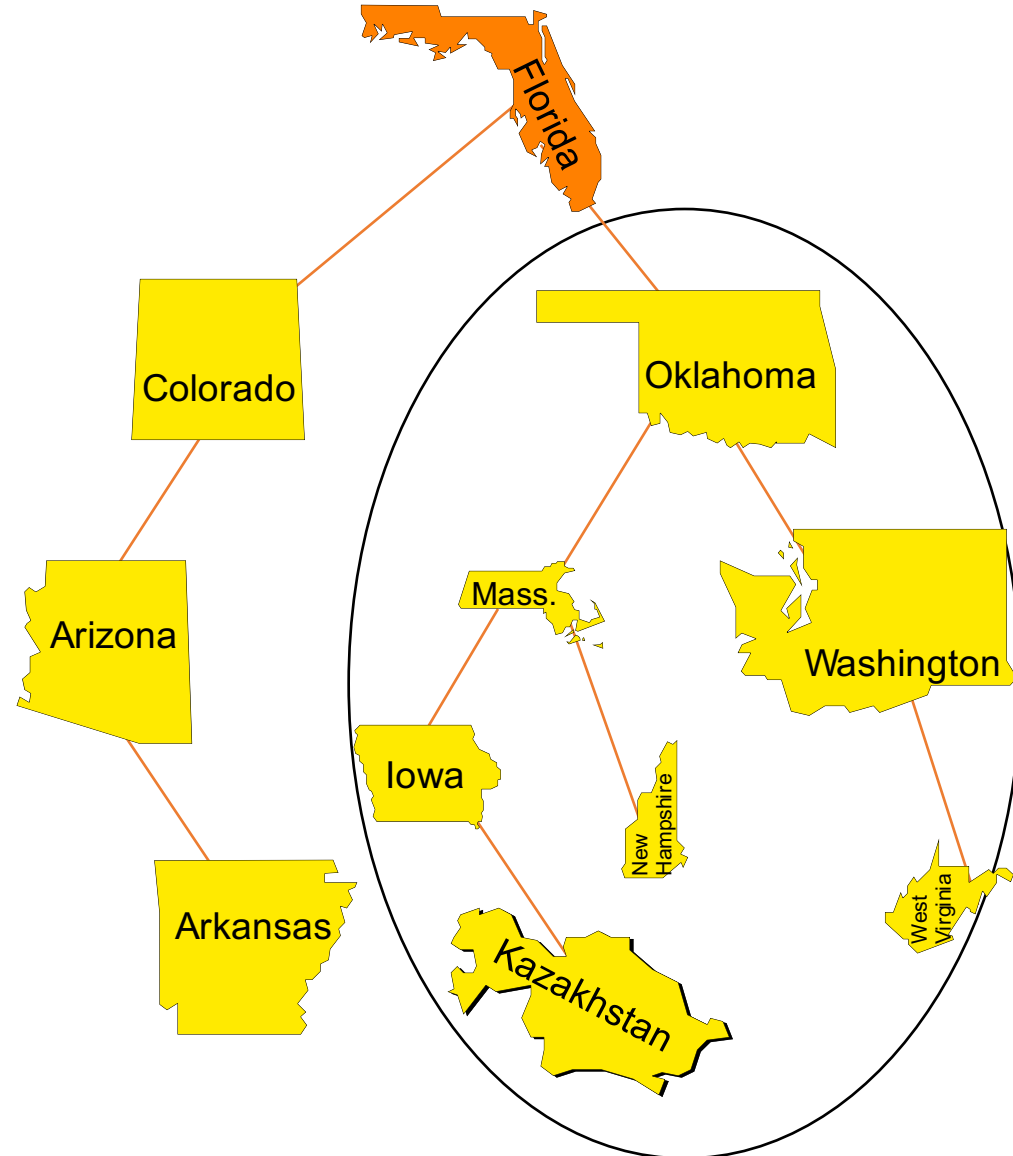
# Removing 'Florida'

The problem of breaking the tree happens because Florida has 2 children.



# Removing 'Florida'

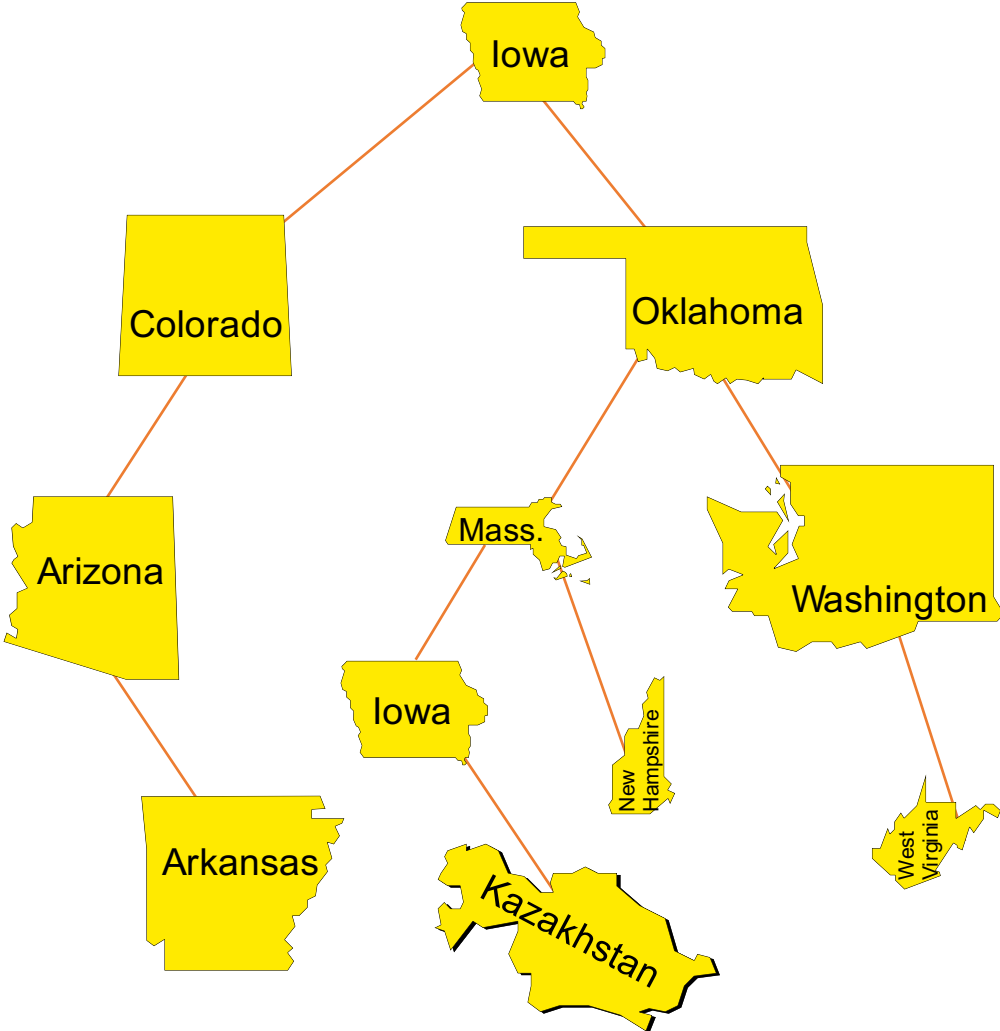
For the rearranging,  
take the **smallest** item  
in the right subtree...



Work for  
multi-set?

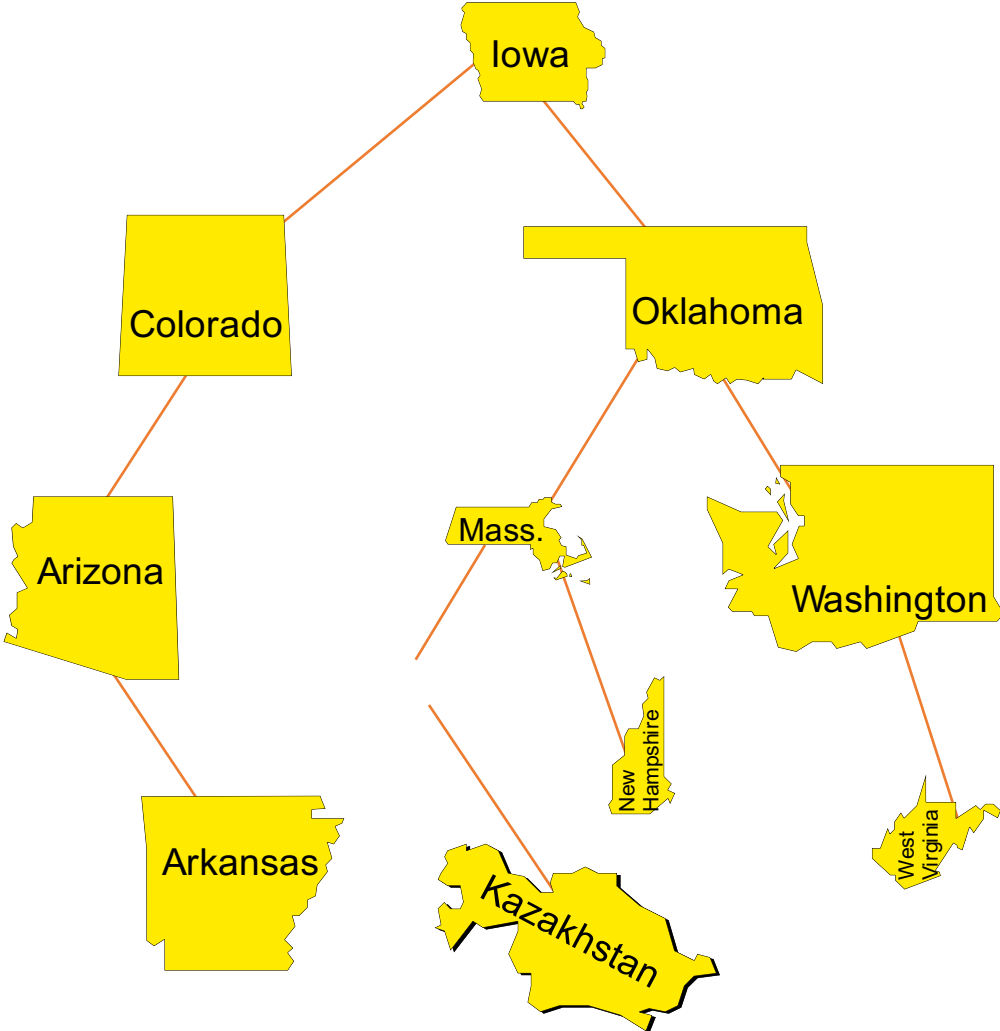
# Removing 'Florida'

...**copy** that smallest item onto the item that we're removing...



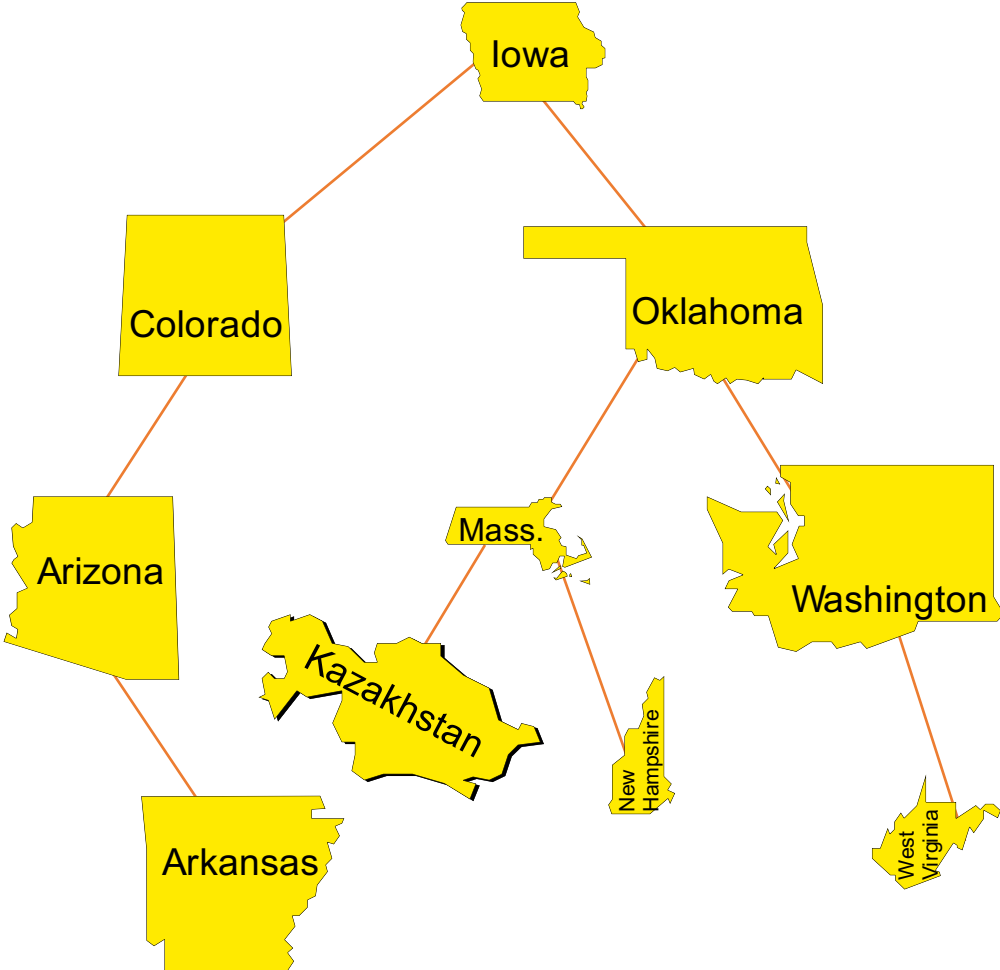
# Removing 'Florida'

... and then remove the extra copy of the item we copied...



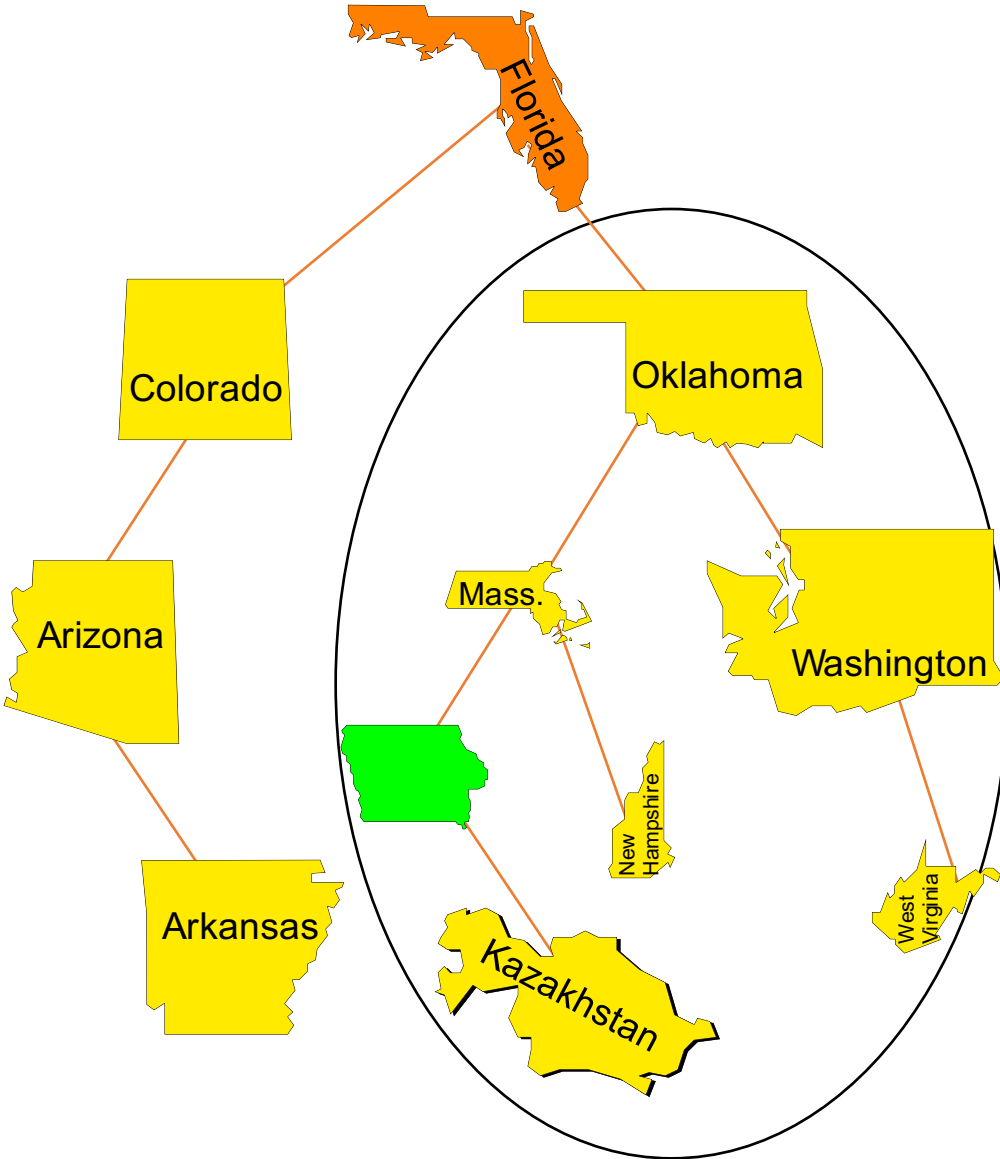
# Removing 'Florida'

... and reconnect the tree



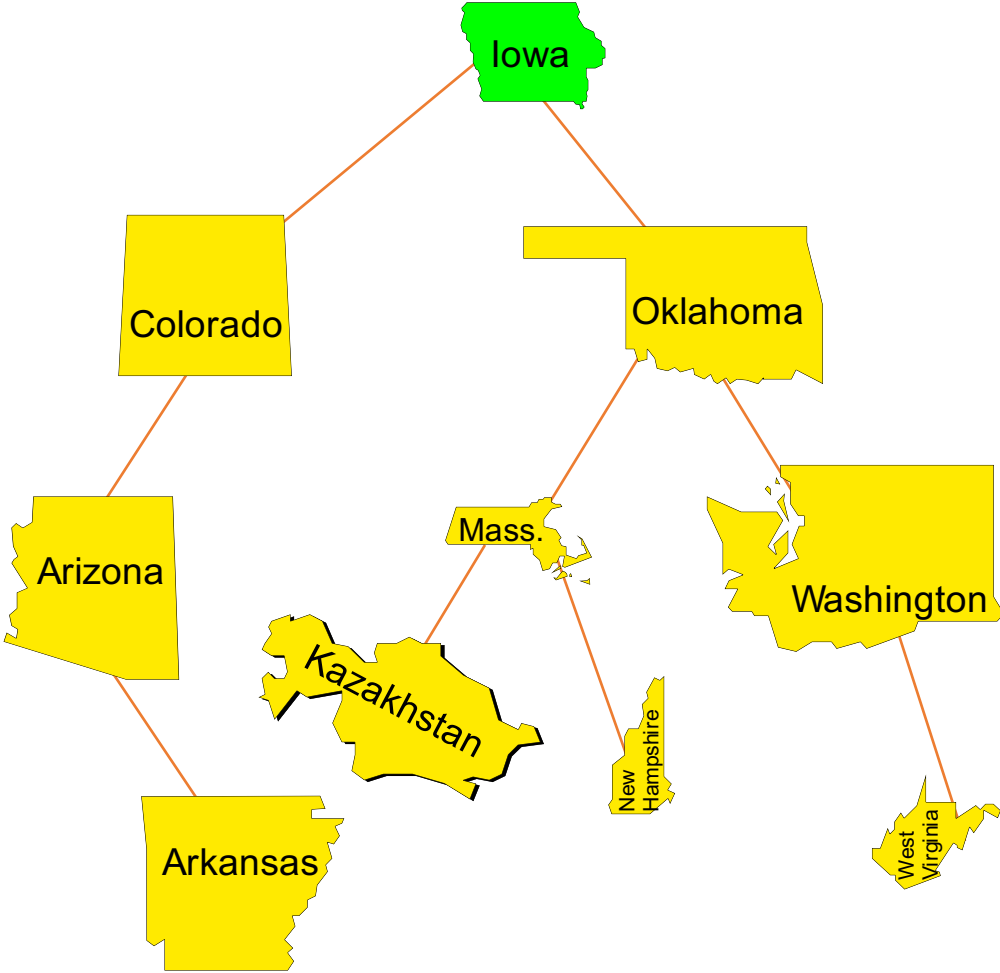
# Removing 'Florida'

*Why did I choose the smallest item in the right subtree?*



# Removing 'Florida'

Because every key must be smaller than the keys in its right subtree





# Removing an Item with a Given Key

- ① Find the item.
  - ② If the item has a right child, rearrange the tree:
    - Find smallest item in the right subtree
    - Copy that smallest item onto the one that you want to remove
    - Remove the extra copy of the smallest item (making sure that you keep the tree connected)
- else just remove the item.



- Binary search trees are a good implementation of data types such as sets, bags, and dictionaries.
- Searching for an item is generally quick since you move from the root to the item, without looking at many other items.
- Adding and deleting items is also quick.
- But as you'll see later, it is possible for the quickness to fail in some cases -- can you see why?

# Assignment

- Read Section 10.5
- Assignment 6 – Bag class with a BST
  - Member functions
    - `void insert(const Item& entry);`
    - `size_type count (const Item& target);`
  - Non-member functions
    - `void bst_remove_all(binary_tree_node<Item>*&root const Item& target);`
    - `void bst_remove_max(binary_tree_node<Item>*&root, Item& removed);`

**Deadline: Monday, November 28, 2016**



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THE END